

HOMOMORPHISMS
STRUCTURAL RAMSEY THEORY
LIMITS

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FoCM MONTEVIDEO
DEC 13, 2014

1

HOMOMORPHISMS STRUCTURAL RAMSEY THEORY LIMITS

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F₀CM MONTEVIDEO
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MODERN DICHOTOMIES

SPARSE

VS

DENSE

STRUCTURE

VS

RANDOMNES

MODEST GOAL

THREE RESULTS

&

ONE CONJECTURE(S)
|
TWO

ALL MATHEMATICAL CONTEXT
OF COMBINATORICS

I

COUNTING HOMOMORPHISM (SUBGRAPHS)

DEF.

A HOMOMORPHISM IS A MAP
WHICH PRESERVES ALL INFORMATION
(DEFINING STRUCTURE)

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FOR GRAPHS:

$f: V \rightarrow V'$ IS HOMOMORPHISM
OF $G=(V,E) \rightarrow G'=(V',E')$
IF $xy \in E \Rightarrow f(x)f(y) \in E'$.

I

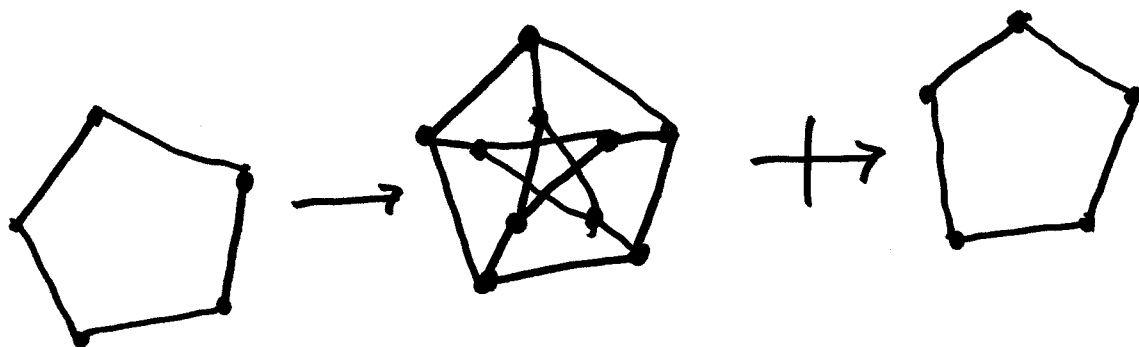
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$$\text{hom}(F, G) = \left| \left\{ f; \underset{\text{HOMO}}{f}: F \rightarrow G \right\} \right|$$

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(~ G-COLOURING OF F)

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~ IMPORTANT INVARIANT

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(\sim G-COLOURING OF F)

\sim PARTITION FUNCTION

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\sim IMPORTANT INVARIANT



WHAT CAN WE SAY ABOUT

$$\frac{\log \text{hom}(F, G)}{\log |G|} \quad ?$$

(i.e. TRYING TO EXPRESS
 $\text{hom}(F, G)$ AS $|G|^\alpha$)

THM (N., P. OSSONA DE MENDEZ)

FOR ANY CLASS \mathcal{C} OF GRAPHS

$$\sup_G \frac{\log \text{hom}(F, G)}{\log |G|} \longrightarrow \{-\infty, 0, 1, \dots, |F|\}$$

FOR ANY F .

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COUNTER EXAMPLE :

$$\mathcal{C} = \{G; G \not\supseteq \square\}$$

SUP NOT INTEGER

$$\left(\frac{3}{2}\right)$$

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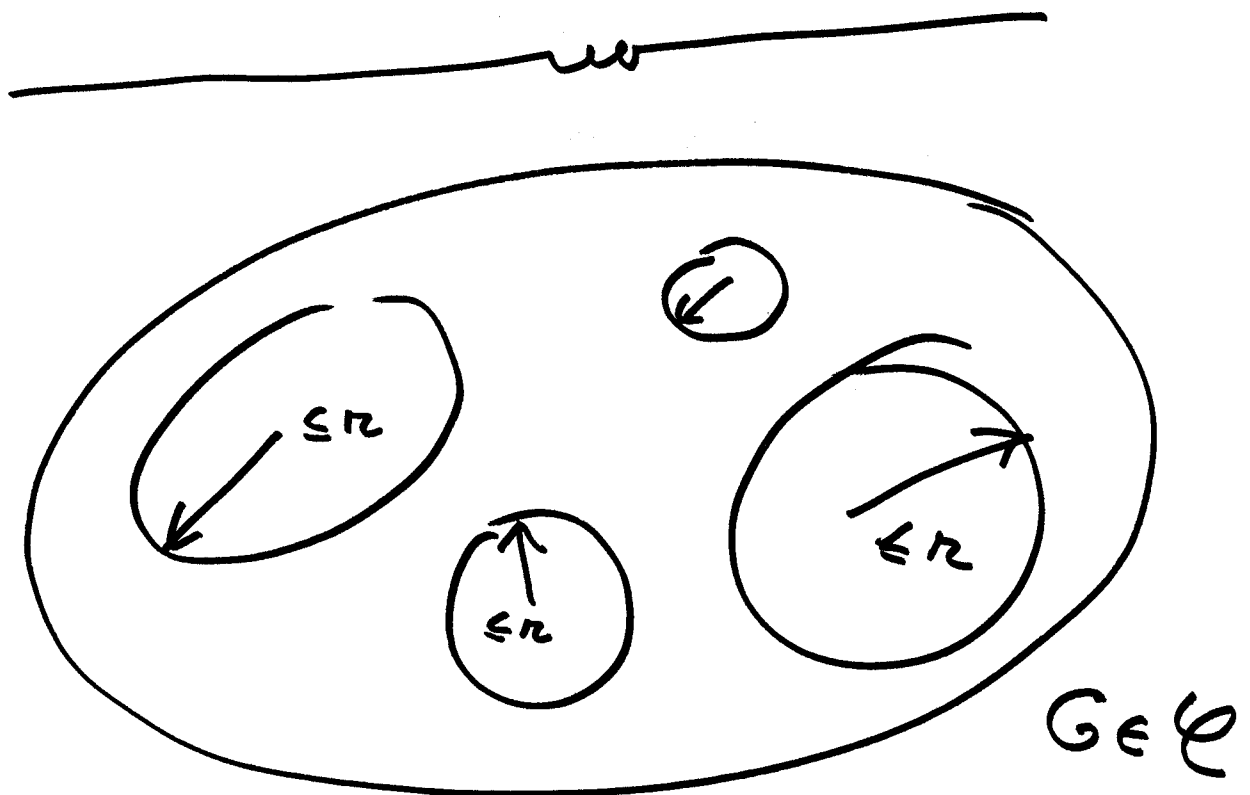
SUP NOT INTEGER $\left(\frac{3}{2}\right)$

CENTRAL QUESTION OF EXTREMAL
THEORY

HIERARCHY OF CLASSES

$$\mathcal{C} \subseteq \mathcal{C}_{\nabla 0} \subseteq \mathcal{C}_{\nabla 1} \subseteq \dots$$

$\mathcal{C}_{\nabla i}$ = ALL GRAPHS WHICH
CAN BE OBTAINED
FROM GRAPHS IN \mathcal{C}
BY CONTRACTING SUBGRAPHS
WITH RADIUS $\leq i$.



THM I

FOR ANY CLASS \mathcal{C} AND ANY F :

$$\lim_{i \rightarrow \infty} \sup_{G \in \mathcal{C} \nabla i} \frac{\log \text{hom}(F, G)}{\log |G|} \in \underbrace{\{-\infty, 0, 1, \dots, |F|\}}_{(\alpha(F))}$$

2

$$\limsup < |F| \Leftrightarrow \mathcal{C} \nabla i \nmid \text{ALL GRAPHS} \\ \text{FOR ALL } i$$

THM I

FOR ANY CLASS \mathcal{C} AND ANY F :

$$\lim_{i \rightarrow \infty} \sup_{G \in \mathcal{C} \nabla i} \frac{\log \text{hom}(F, G)}{\log |G|} \in \underbrace{\{-\infty, 0, 1, \dots, |F|\}}_{\alpha(F)}$$

$\&$

$$\limsup < |F| \Leftrightarrow \mathcal{C} \nabla i \nmid \text{ALL GRAPHS}$$

FOR ALL i

(NOWHERE DENSE CLASSES)

VS

SOMEWHERE DENSE CLASSES

(DVORÁK, NORINE, KREUTZER, GROHE,
KRAJČAL, THOMAS, ...)

Algorithms and Combinatorics 28

Jaroslav Nešetřil

Patrice Ossona de Mendez

Sparsity

Graphs, Structures, and Algorithms



Springer

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COROLLARY

\mathcal{C} NOWHERE DENSE



$$\lim_{i \rightarrow \infty} \limsup_{G \in \mathcal{C} \cap \mathcal{C}_i} \frac{\log(\#F \subseteq G)}{\log |G|}$$

\in

$$\{-\infty, 0, 1, \dots, \alpha(F)\}$$

FOR EVERY F .

WITH AT LEAST
ONE EDGE
SO THAT $\alpha(F) < |F|$.

PROOF

IT SUFFICES TO PROVE



PROOF COMBINES ALL CHARACTERIZATION

— TREE DEPTH td

— LOW TREE DEPTH DECOMPOSITIONS
 χ_p

— STEPPING UP LEMMA

FOR " (k, F) -SUNFLOWERS"

— COUNTING OF COLORED TREES

PROOF II

FIRST CASE: $F = K_2$

THM

(TRICHOTOMY THM)

FOR EVERY INFINITE CLASS \mathcal{C}
THE LIMIT

$$\lim_{n \rightarrow \infty} \limsup_{G \in \mathcal{C} \nabla i} \frac{\log |G|}{\log |G|}$$

$$\wedge \{ -\infty, 0, 1, 2 \}$$

$$\leq 1$$

IFF

\mathcal{C} IS NOWHERE DENSE

$$2$$

IFF

\mathcal{C} IS SOMEWHERE DENSE

II

CHARACTERISATION OF RAMSEY CLASSES

FINITE RAMSEY THEOREM

$$\forall p, k, n \exists N = N(p, k, n) : N \rightarrow (n)_k^p$$

$$\text{IF } |X| = N, \binom{X}{p} = \{A \subseteq X; |A| = p\}$$

$$\binom{X}{p} = a_1 \cup \dots \cup a_k \quad (\text{ARBITRARY PARTITION})$$

THEN

$$\exists Y, Y \subseteq X, |Y| = n \exists i_0$$

$$\binom{Y}{p} \subseteq a_{i_0}$$


DEF RAMSEY CLASS \mathcal{K}
(OF STRUCTURES)

$$A, B, C, \dots \in \mathcal{K}$$

$\binom{B}{A}$ ALL SUBSTRUCTURES OF B
ISOMORPHIC TO A

\mathcal{K} IS RAMSEY CLASS (LEEB,
N., RODL)
IFF

$$\forall A \forall B \forall k \exists C : C \rightarrow \binom{B}{k}^A$$



IN \mathcal{K}

(ERDŐS-RADO) PARTITION ARROW :

FOR EVERY $\binom{C}{A} = a_1 \cup \dots \cup a_k \quad \exists B' \in \binom{C}{B}$

SUCH THAT $\binom{B'}{A} \subseteq a_{i_0}$

RAMSEY CLASSES \equiv TOP OF THE LINE
OF "RAMSEY
PROPERTIES"

EXAMPLES

ALL COMPLETE GRAPHS (FRT)
ALL FINITE LINEAR ORDERS (FRT)
ALL FINITE VECTOR SPACES
ALL COMBINATORIAL CUBES

NOT
GRAPHS,

NECESSARY CONDITIONS FOR RAMSEY CLASSES

①

RIGIDITY

(ORDERING, ORDERING LEMMA)

N. RODL

ANGEL, KECHRIS, LYONS

NECESSARY CONDITIONS FOR RAMSEY CLASSES

① RIGIDITY

(ORDERING, ORDERING LEMMA)

N. RODL

ANGEL, KECHRIS, LYONS

② GLOBAL SYMMETRY

(AGE OF ULTRAHOMOGENEOUS
STRUCTURE)

N.

KECHRIS, PESTOV, TODORCEVIC

CHARACTERIZATION PROGRAM

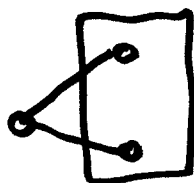
HOMOMORPHISM DEFINE CLASSES

$$\mathcal{F} = \{F_1, F_2, \dots, F_t\}$$

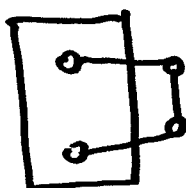
$$\text{FORB}(\mathcal{F}) = \left\{ G; \begin{array}{l} F_i \not\rightarrow G \\ i=1, \dots, t \end{array} \right\}$$

$\text{FORB}(\square) = \text{ALL GRAPHS NOT CONTAINING } \triangle \text{ OR } \square$

$\text{FORB}(\square) \text{ NOT RAMSEY}$



BLUE



RED

BUT

$\text{FORB}(\square) + \text{BLUE-RED}$



LIFT OF $\text{FORB}(\square)$
(EXPANSION)

IS RAMSEY

BUT

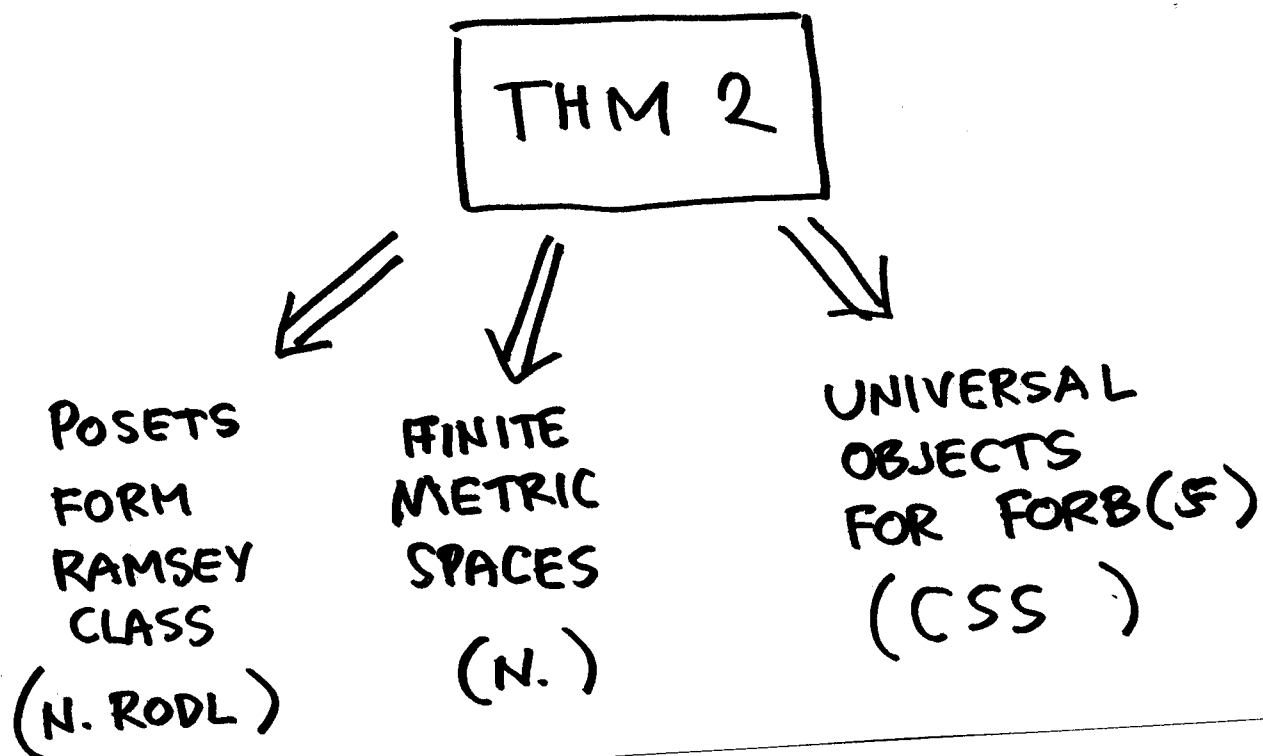
$\text{FORB}(\square) + \text{BLUE-RED}$

LIFT OF $\text{FORB}(\square)$
(EXPANSION)

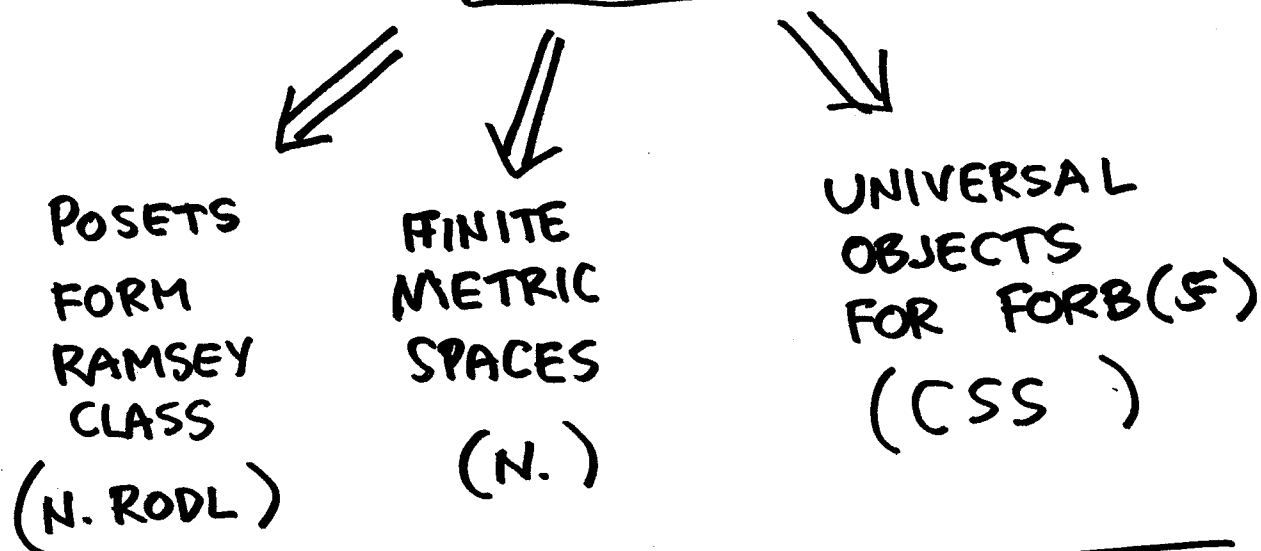
IS RAMSEY

THM II (N. 13)

FOR EVERY FINITE SET OF
CONNECTED GRAPHS \mathcal{F}
THE CLASS $\text{FORB}(\mathcal{F})$ HAS
A FINITE LIFT WHICH FORM RAMSEY
CLASS.



THM 2

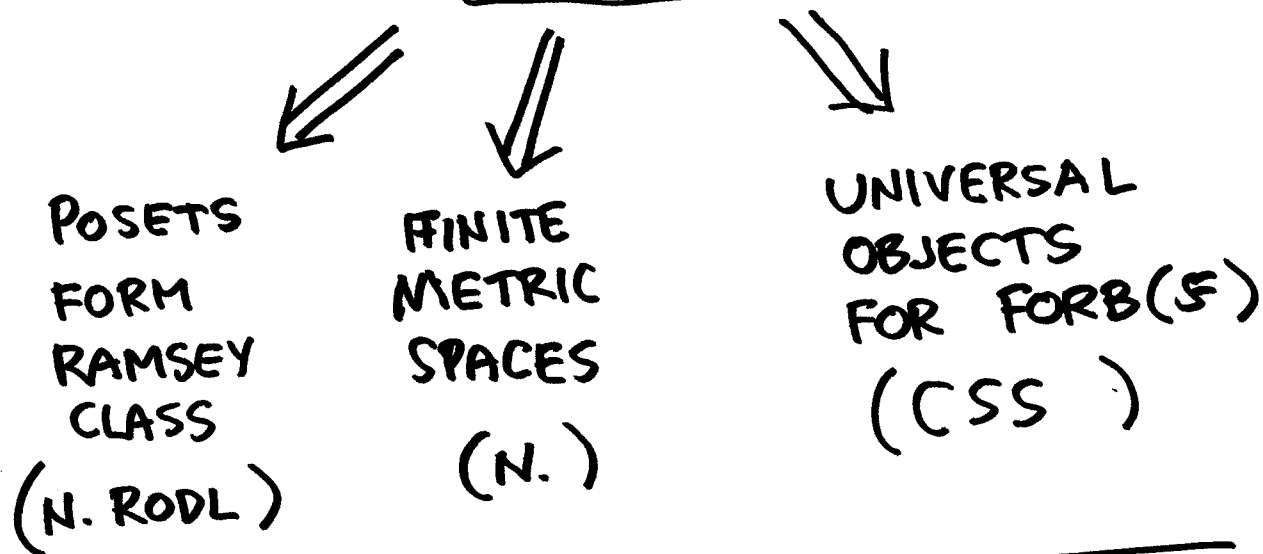


PROBLEM

EVERY ULTRAHOMOGENEOUS STRUCTURE HAS

A FINITARY LIFT WHICH FORM RAMSEY CLASS.

THM 2



PROBLEM

EVERY ULTRAHOMOGENEOUS
STRUCTURE HAS

A FINITARY LIFT WHICH FORM
RAMSEY CLASS.

BOWTIE-FREE GRAPHS (KOMJATH)
HAVE RAMSEY LIFT
(HUBIČKA, N.)

COROLLARY

EVERY HOMOMORPHISM CLOSED
FIRST ORDER DEFINABLE CLASS
OF STRUCTURES HAS A RAMSEY
LIFT.

(PF : THEOREM II
+
ROSSMAN)

III. LIMITS BY SATISFACTION

$$\exists f: F \longrightarrow G$$



$$G \models \varphi_f$$

$$\left(\varphi_f = \bigwedge_{ij \in E(F)} x_i x_j \right)$$

φ FIRST ORDER FORMULA

$\varphi(x_1, \dots, x_t)$ WITH t FREE VARIABLES

$$\varphi(G) = \{ (v_1, \dots, v_t); G \models \varphi(v_1, \dots, v_t) \}$$

$$\langle \varphi, G \rangle = \frac{|\varphi(G)|}{|G|^t}$$

φ FIRST ORDER FORMULA

$\varphi(x_1, \dots, x_t)$ WITH t FREE VARIABLES

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$$\langle \varphi, G \rangle = \frac{|\varphi(G)|}{|G|^t}$$

"STONE BRACKET"

PROBABILITY THAT t VERTICES
CAN BE EXTENDED TO SATISFACTION
OF φ

EXAMPLE

$$\varphi(x_1, x_2) = \exists x \left(xx_1 \in E \text{ \& } xx_2 \in E \right) \\ \& \\ x_1 \neq x_2$$

|||

BETWEEN x_1 AND x_2 THERE IS A PATH OF LENGTH 2.

$\langle \varphi, G \rangle$ = PROB. THAT 2 VERTICES OF G ARE JOINED BY A PATH OF LENGTH 2.

$G_1, G_2, \dots, G_n, \dots$

IS

FO -
CONVERGENT

IFF

FOR EVERY φ THE NUMBERS

$\langle \varphi, G_i \rangle$

CONVERGE .

N. POSSONA
DE HENDEZ

$G_1, G_2, \dots, G_n, \dots$

IS

| |
|--------------------|
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|--------------------|

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CONVERGE.

N. POSSONA
DE HENDEZ

X - CONVERGENCE

⋮

FOR EVERY $\varphi \in X$

— — — —

$\mathcal{F}_0(X)$ \cup X

BOOLEAN ALGEBRA

BOOLEAN SUBALGEBRA

 $S(X)$ STONE SPACE OF X (ULTRAFILTERS OR
HOMOMORPHISMS $f: X \rightarrow \{0,1\}$)

+ TOPOLOGY BY

$$K_X(\varphi) = \{U \mid \varphi \in U\}$$

$$= \{f \mid f(\varphi) = 1\}$$

STONE DUALITY

THM

X SUBALGEBRA OF $FO(\mathcal{L})$.

FOR EVERY X -CONVERGENT SEQUENCE $(G_n)_{n \in \mathbb{N}}$ THERE EXISTS UNIQUE PROBABILITY MEASURE μ ON $S(X)$ SUCH THAT FOR EVERY FORMULA $\varphi(x_1, \dots, x_p) \in X$ HOLDS:

$$\int_{S(X)} \mathbb{1}_{K(\varphi)}(x) d\mu(x) =$$

$$= \lim_{n \rightarrow \infty} \langle \varphi, G_n \rangle$$

EXTREME CASE

\mathbb{G} MODELLING IS STANDARD
BOREL SPACE

SUCH THAT EVERY FIRST ORDER
DEFINABLE SET IS MEASURABLE

EXTREME CASE

\mathbb{G} MODELLING IS STANDARD
BOREL SPACE

SUCH THAT EVERY FIRST ORDER
DEFINABLE SET IS MEASURABLE

THM 3

EVERY ^{FO} CONVERGENT SEQUENCE (G_n)
CONVERGES TO MODELLING



\mathcal{C} IS NOWHERE DENSE

CONVERSE TRUE FOR

- GRAPHS WITH BOUNDED TREE DEPTH
- GRAPHS WITH BOUNDED DEGREES
- TREES



?



?

