# O teorema de Gallai sobre caminhos disjuntos

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O teorema de Gallai<sup>1</sup> é uma generalização natural do teorema dos emparelhamentos máximos (Teorema 2.2.3 em Diestel<sup>2</sup>). Ele é a base de uma nova demonstração do teorema de Mader (Teorema 3.4.1 em Diestel) recentemente descoberta por Alex Schrijver (A short proof of Mader's S-paths theorem)<sup>3</sup>.

Estou tratando do assunto como se ele fosse a seção 2.4 do capítulo 2 do livro de Diestel (se bem que a seção deveria estar entre a 2.2 e a 2.3). A propósito, seria muito conveniente que Diestel substituisse a discussão no final da seção 2.2 (página 38) pelo seguinte corolário do Teorema 2.2.3:

Corollary 2.2.4 Every graph G has a matching M and a disjoint collection  $\mathcal{Y}$  of subsets of V(G) such that no edge links two different members of  $\mathcal{Y}$  and  $|M| = |V(G) \setminus \mathcal{Y}| + \sum_{Y \in \mathcal{Y}} \lfloor \frac{1}{2} |Y| \rfloor$ .

## 2.4 Gallai's theorem

Let A be a set of vertices of a graph G; we shall regard A as fixed throughout this section. An A-path is any path<sup>4</sup> whose ends are two distinct vertices in A and whose internal vertices are in  $V(G) \setminus A$ .<sup>5</sup>

disjoint collection

A-path

A collection  $\mathcal{P}$  of A-paths is *disjoint* if no vertex of the graph belongs to more than one member of  $\mathcal{P}$ . A matching, for example, is essentially the same thing as a disjoint collection of V(G)-paths.

split width of a split

Our problem in this section is the characterization of maximum disjoint collections of A-paths. This will be done in terms of splits<sup>6</sup>. A *split* of G is just a disjoint collection of subsets of V(G). The A-width of a split  $\mathcal{Y}$  is the number

$$\operatorname{wd}_A(\mathcal{Y}) := |\overline{\bigcup \mathcal{Y}}| + \sum_{Y \in \mathcal{Y}} \left\lfloor \frac{|Y \cap A|}{2} \right
floor,$$

<sup>&</sup>lt;sup>1</sup> Tibor Gallai, Maximum-minimum Sätze und verallgemeinerte Faktoren von Graphen, Acta Mathematica Academiae Scientiarum Hungaricae 12 (1961), 131–173

<sup>&</sup>lt;sup>2</sup> Graph Theory, 2nd. edition, Springer, 1999

<sup>&</sup>lt;sup>3</sup> tomei conhecimento por intermédio de José Coelho de Pina Jr.

<sup>&</sup>lt;sup>4</sup> by definition, paths have no repreated vertices

<sup>&</sup>lt;sup> $^{5}$ </sup> we could drop the requirement that no internal vertex be in A, but this would have no practical consequence

<sup>&</sup>lt;sup>6</sup> this term is not standard

where  $\overline{U}$  denotes  $V(G) \setminus U$ . A split  $\mathcal{Y}$  is *strong* if no edge of G links two different—strong split members of  $\mathcal{Y}$ .

**Lemma 2.4.1** For any disjoint collection  $\mathcal{P}$  of A-paths and any strong split  $\mathcal{Y}$ ,

$$|\mathcal{P}| \leq \operatorname{wd}_A(\mathcal{Y})$$
.

Proof. Let  $X := \overline{\bigcup \mathcal{Y}}$ . Each A-path has a vertex in X or at least two vertices (its terminal vertices) in  $Y \cap A$  for some Y in  $\mathcal{Y}$ . With each P in  $\mathcal{P}$ , we may associate (1) a vertex in  $V(P) \cap X$  if this set is not empty or (2) a pair of vertices in  $A \cap Y$ , where Y is the appropriate member of  $\mathcal{Y}$ . Each vertex in X will have been associated with at most one member of  $\mathcal{P}$ , since  $\mathcal{P}$  is disjoint. On the other hand, in each set of the form  $Y \cap A$  a disjoint collection of pairs of vertices will have been defined; each of these pairs will have been associated with a different member of  $\mathcal{P}$ , since  $\mathcal{P}$  is disjoint. The inequality  $|\mathcal{P}| \leq \operatorname{wd}_A(\mathcal{Y})$  follows.

The case A = V(G) of the lemma has been studied in section 2.2: it amounts to the inequality  $|M| \leq \operatorname{wd}_{V(G)}(\mathcal{Y})$ , valid for any matching M and any strong split  $\mathcal{Y}$ .

**Theorem 2.4.2** (Gallai 1961) There exists a disjoint collection  $\mathcal{P}$  of A-paths and a strong split  $\mathcal{Y}$  such that

[3.4.1]

E

$$|\mathcal{P}| = \operatorname{wd}_A(\mathcal{Y})$$
.

In other words, the A-width of a minimum-width strong split is equal to the maximum number of disjoint A-paths. This is a generalization of Theorem 2.2.3. That theorem asserts there exists a matching M and a strong split  $\mathcal{Y}$  such that  $|M| = \operatorname{wd}_{V(G)}(\mathcal{Y})$ .

**Proof of Theorem 2.4.2.** Let G' and G'' be two mutually disjoint copies of G. For each vertex v of G, let v' and v'' denote the corresponding vertices in G' and G'' respectively. Let A' and A'' be the copies of A in G' and G''. Let V := V(G), A', A'' V' := V(G'), and V'' := V(G'').

Let E be the set of all pairs v'v'' with v a vertex of G and let C the set of all pairs of the form v'w'' such that vw is an edge of G. Now define

$$H := ((G' \cup G'') + E + C) - A''.$$

Note that for each edge vw of G the vertices v', w', v'', and w'' are pairwise adjacent in H. If v'w' is an edge of H then so are v'w'' and v''w' (unless  $w'' \in A''$  or  $v'' \in A''$ ), and if v'w'' is an edge of H with  $v \neq w$  then also v'w' and v''w'' are edges of H (unless  $v'' \in A''$ ).

Let  $\mathcal{Y}$  be a minimum-width strong split of H. Theorem 2.2.3 guarantees that (2.2.3)

$$|M| = \operatorname{wd}_{V(H)}(\mathcal{Y})$$

for some matching M in H. Choose the split  $\mathcal{Y}$ , among all those having minimum  $M, \mathcal{Y}$ 

width, so that  $\bigcup \mathcal{Y}$  is maximal. Then, for each v in  $V \setminus A$ ,

vertices 
$$v'$$
 and  $v''$  are either both in  $X$  or both not in  $X$ , (2.1)

where  $X := \overline{\bigcup \mathcal{Y}}$ . In order to prove this claim, suppose for a while that there exists a vertex v in  $V \setminus A$  such that  $v' \in X$  but  $v'' \notin X$ . Then  $v'' \in Y$  for some Y in  $\mathcal{Y}$ . Now let  $\tilde{Y} := Y \cup \{v'\}$  and  $\tilde{\mathcal{Y}} := (\mathcal{Y} \setminus \{Y\}) \cup \{\tilde{Y}\}$ . Observe that

$$\operatorname{wd}_{V(H)}(\tilde{\mathcal{Y}}) \leq \operatorname{wd}_{V(H)}(\mathcal{Y})$$
.

Hence, the maximality of  $\bigcup \mathcal{Y}$  implies that the split  $\tilde{\mathcal{Y}}$  is not strong, i.e., H has an edge v'w' with w' in  $V' \setminus (X \cup Y)$  or an edge v'w'' with w'' in  $V'' \setminus (X \cup Y)$ . In the first case, v''w' is an edge of H and it connects Y to some other member of  $\mathcal{Y}$ , which is impossible because  $\mathcal{Y}$  is strong. In the second case, v''w'' is an edge of H and it connects Y to some other member of  $\mathcal{Y}$ , which is again impossible. This contradiction proves claim (2.1) in case  $v' \in X$ . The proof in case  $v'' \in X$  is similar.

Let 
$$X' := X \cap V'$$
 and  $X'' := X \cap V''$ . By virtue of (2.1),  $|X' \setminus A'| = |X''|$ . Hence  $X', X'' \mid X \mid = |X' \cap A'| + 2|X' \setminus A'|$ .

For each Y in  $\mathcal{Y}$ , let  $Y':=Y\cap V'$  and  $Y'':=Y\cap V''$ . For each v in  $V\setminus A$ , by virtue of (2.1), either Y contains both v' and v'' or none of these two vertices. Hence,  $|Y'\setminus A'|=|Y''|$ , and therefore  $|Y|=|Y'\cap A'|+2|Y'\setminus A'|$ . The collection  $\mathcal{Y}'$  of all Y' is a strong split of G' and

$$\operatorname{wd}_{A'}(\mathcal{Y}') = \operatorname{wd}_{V(H)}(\mathcal{Y}) - |V' \setminus A'|.$$

Here is the proof of this equality, with k denoting the number of Y in  $\mathcal{Y}$  for which |Y| is odd, and therefore also the number of Y' in  $\mathcal{Y}'$  for which  $|Y' \cap A'|$  is odd:

$$\begin{aligned}
\operatorname{wd}_{V(H)}(\mathcal{Y}) &= |X| - \frac{k}{2} + \sum_{Y \in \mathcal{Y}} \frac{|Y|}{2} \\
&= |X' \cap A'| + 2|X' \setminus A'| - \frac{k}{2} + \sum_{Y' \in \mathcal{Y}'} \frac{|Y' \cap A'| + 2|Y' \setminus A'|}{2} \\
&= |X'| + |X' \setminus A'| - \frac{k}{2} + \sum_{Y'} \frac{|Y' \cap A'|}{2} + \sum_{Y'} |Y' \setminus A'| \\
&= |X'| + \sum_{Y'} \left\lfloor \frac{|Y' \cap A'|}{2} \right\rfloor + |V' \setminus A'| \\
&= \operatorname{wd}_{A'}(\mathcal{Y}') + |V' \setminus A'| ,
\end{aligned}$$

Now we must convert matching M into a large disjoint collection of A'-paths. Let E' be the set of edges in E incident to  $V' \setminus A'$ ; of course E' is a matching in H with exactly  $|V' \setminus A'|$  edges. Since  $|M| = \operatorname{wd}_{V(H)}(\mathcal{Y}) = \operatorname{wd}_{A'}(\mathcal{Y}') + |E'|$ , the graph  $H[E' \cup M]$  has at least  $\operatorname{wd}_{A'}(\mathcal{Y}')$  components with more edges in M than in E'. Each such component is a path whose first and last edges are in M. The ends of such path must be both in A'; this is only possible if its first and last edges have both ends in

V' and therefore each of its edges in M has both ends in V' or both in V''. Hence, each of these paths has the form

$$v'_0 \ v'_1 \ v''_1 \ v''_2 \ v'_2 \ v''_3 \ v''_3 \ v''_4 \ \dots \ v''_{k-1} \ v''_k \ v'_k \ v'_{k+1} \ ,$$

where  $v_0$  and  $v_{k+1}$  are in A and  $v_1, \ldots, v_k$  are in  $V \setminus A$ . Of course

$$v_0' v_1' v_2' v_3' v_4' \dots v_k' v_{k+1}'$$

is an A'-path in G'. Moreover, such paths are pairwise disjoint. So, we have a disjoint collection of

$$\operatorname{wd}_{A'}(\mathcal{Y}')$$

A'-paths in G'. Since G' is isomorphic to G, our result follows.

An immediate consequence of Lemma 2.4.1 and Theorem 2.4.2 is the minimax identity  $\max |\mathcal{P}| = \min \operatorname{wd}_A(\mathcal{Y})$ , where  $\max$  is taken over all disjoint collections  $\mathcal{P}$  of A-paths and  $\min$  is taken over all strong splits  $\mathcal{Y}$ .

### Exercises

1. Let G' and G'' be two mutually disjoint copies of a graph G. For each vertex v of G, let v' and v'' denote the corresponding vertices in G' and G'' respectively. Let E denote the set of all pairs of the form v'v'' with v in V(G). Let C denote the set of all pairs of the form v'w'' such that vw is an edge of G. Finally, let

$$H := (G' \cup G'') + E + C$$
.

Describe, in detail, the components of the graph  $H[M \cup E]$ , where M is a matching in H. Now suppose one of the components of  $H[M \cup E]$  is a path of the form

$$v'_1 v''_1 v''_2 v''_2 v''_3 v''_3 v''_4 v'_4 \dots v'_{k-1} v'_k v''_k$$

and show that  $v'_1 \ v'_2 \ v'_3 \ v'_4 \ \dots \ v'_{k-1} \ v'_k$  is a path in G'.

- 2. Describe an algorithm for deciding whether a split  $\mathcal{Y}$  is strong.
- 3. From Lemma 2.4.1, show that  $|\mathcal{P}| \leq \lfloor \frac{1}{2}|A| \rfloor$  for any disjoint collection  $\mathcal{P}$  of A-paths.
- 4. Give the terms  $(\mathcal{P} \text{ and } \mathcal{Y})$  of Gallai's theorem in case  $|A| \leq 1$ . Repeat for |A| = 2. Repeat for |A| = 3.
- 5. Complete the proof of claim (2.1), i.e., show that there is no v in  $V \setminus A$  such that  $v' \notin X$  but  $v'' \in X$ .

6. Show that Theorem 2.4.2 could have been stated as follows: There exists a disjoint collection  $\mathcal{P}$  of A-paths and a set X of vertices such that

$$|\mathcal{P}| = |X| + \sum_{C} \left\lfloor \frac{|V(C) \cap A|}{2} \right
floor,$$

where the sum is taken over all components C of G-X.