Character-Theoretic Tools for Studying Linear Codes over Rings and Modules

Jay A. Wood

Department of Mathematics Western Michigan University http://sites.google.com/a/wmich.edu/jaywood

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5. Exercise Session

- Exercise: every character of Z/kZ has the form ρ_b(a) = exp(2πiab/k), a ∈ Z/kZ, for some b ∈ Z/kZ. [What is ρ(1)?]
- Thus, $(\mathbb{Z}/k\mathbb{Z})^{\widehat{}} \cong \mathbb{Z}/k\mathbb{Z}$, via $\rho_b \longleftrightarrow b$.

Duality functor

- Pontryagin duality: $A \mapsto \widehat{A}$
- Exact contravariant functor:

$$0 \rightarrow A_1 \rightarrow A_2 \rightarrow A_3 \rightarrow 0$$

induces

$$0
ightarrow \widehat{A}_3
ightarrow \widehat{A}_2
ightarrow \widehat{A}_1
ightarrow 0.$$

≅ A, but not naturally. (*Uses fundamental theorem of finitely generated abelian groups.*)
 ≅ A, naturally: a ↦ (π ↦ π(a)).
(A × B) ≃ Â × B̂.

Annihilators

Let B ⊆ A be any subgroup.
Define the annihilator (Â : B):

$$(\widehat{A}:B) = \{ \rho \in \widehat{A}: \rho(B) = 1 \} = \{ \varrho \in \widehat{A}: \varrho(B) = 0 \}.$$

- $(\widehat{A}:B)\cong (A/B)^{\widehat{}}.$
- $|B| \cdot |(\widehat{A} : B)| = |A|.$
- Double annihilator: $(A : (\widehat{A} : B)) = B$.

Summation formulas

- Need multiplicative form of characters.
- For $\pi \in \widehat{A}$,

$$\sum_{a\in A}\pi(a)=egin{cases} |A|,&\pi=1,\ 0,&\pi
eq 1. \end{cases}$$

• For $a \in A$,

$$\sum_{\pi\in\widehat{A}}\pi(a)=egin{cases} |A|,&a=0,\ 0,&a
eq 0. \end{cases}$$

Fourier transform

• Given a function $f : A \to V$, V a complex vector space. Define its **Fourier transform** $\hat{f} : \hat{A} \to V$ by

$$\widehat{f}(\pi) = \sum_{a \in A} \pi(a) f(a), \quad \pi \in \widehat{A}.$$

•
$$\widehat{}: F(A, V) \to F(\widehat{A}, V).$$

Invert:

$$f(a)=rac{1}{|\mathcal{A}|}\sum_{\pi\in\widehat{\mathcal{A}}}\pi(-a)\widehat{f}(\pi), \quad a\in\mathcal{A}.$$

Poisson summation formula

Let *B* be any subgroup of *A*, and let $f : A \rightarrow V$. Then for any $a \in A$,

$$\sum_{b\in B}f(a+b)=rac{1}{|(\widehat{A}:B)|}\sum_{\pi\in(\widehat{A}:B)}\pi(-a)\widehat{f}(\pi).$$

If a = 0, then

$$\sum_{b\in B} f(b) = rac{1}{|(\widehat{A}:B)|} \sum_{\pi\in(\widehat{A}:B)} \widehat{f}(\pi).$$

A Fourier transform example

- ► Suppose *V* is a complex algebra.
- Suppose $f : A^n \to V$ has the form

$$f(a_1,\ldots,a_n)=\prod_{i=1}^n f_i(a_i),$$

where $f_i : A \rightarrow V$. Then

$$\hat{f}(\pi_1,\ldots,\pi_n)=\prod_{i=1}^n \hat{f}_i(\pi_i).$$

Character modules

- Extra information: the left *R*-module structure on *A* induces a right *R*-module structure on *Â*.
- For $r \in R$ and $\varpi \in \widehat{A}$, define $\varpi r \in \widehat{A}$ by $(\varpi r)(a) = \varpi(ra)$, $a \in A$; $(\pi^r)(a) = \pi(ra)$.
- If A is a right module, then A is a left module:
 (r∞)(a) = ∞(ar); (^rπ)(a) = π(ar).

Annihilators are submodules

- Suppose $B \subseteq A$ is a left *R*-submodule.
- Then the annihilator (Â : B) ⊆ Â is a right R-submodule.
- If $\varrho \in (\widehat{A} : B)$ and $r \in R$, then

$$(\varrho r)(B) = \varrho(rB) \subseteq \varrho(B) = 0,$$

because B is a left submodule.

Characterizing generating characters

Theorem

A character $\varrho \in \widehat{R}$ is a right generating character if and only if ker ϱ contains no nonzero right ideal of R.

- Define $\psi : R \to \widehat{R}$ by $\psi(r) = \varrho r$. When is ψ an isomorphism? (Injective is enough, as $|R| = |\widehat{R}|$.)
- ► $\psi(r) = 0$ iff $(\varrho r)(R) = 0$ iff $\varrho(rR) = 0$ iff $rR \subseteq \ker \varrho$.
- Similar result for left generating characters.

Left/right symmetry

Theorem

A character $\varrho \in \widehat{R}$ is a left generating character if and only if ϱ is a right generating character.

- Left implies right: Suppose $rR \subseteq \ker \varrho$. Then $\varrho(rs) = 0$ for all $s \in R$.
- Then $(s\varrho)(r) = 0$ for all $s \in R$. I.e., $\varpi(r) = 0$ for all $\varpi \in \widehat{R}$, as ϱ left generates.

• Thus r = 0. (Uses " $|B| \cdot |(\widehat{A} : B)| = |\widehat{A}|$ ", $B = \mathbb{Z}r$.)

A generalization for modules

- ► *R* finite ring with 1; *A* finite unital left *R*-module.
- An *R*-module is cyclic if it is generated by one element. Say *M* is generated by *m* ∈ *M*. Then *R* → *M*, *r* → *rm*, is onto.

Theorem

The following are equivalent:

- 1. \widehat{A} is a cyclic right R-module.
- 2. A injects into \widehat{R} : $A \hookrightarrow \widehat{R}$.
- 3. There exists $\varrho \in \widehat{A}$ such that ker ϱ contains no nonzero left *R*-submodule.

Proof

- ▶ 1 ↔ 2. Contravariant exact functor: $0 \to A \to \widehat{R}$ dualizes to $R \to \widehat{A} \to 0$, and vice versa.
- Fix $\rho \in \widehat{A}$. Define $A \to \widehat{R}$ by $a \mapsto (r \mapsto \rho(ra))$.
- 2 ↔ 3: a ∈ A is in the kernel of the map above iff *ρ*(Ra) = 0 iff Ra ⊆ ker *ρ*.
- Call such a *Q* a generating character for *A*.

More on simple modules

- If S is simple, and $0 \neq s \in S$, then S = Rs.
- The annihilator ann(s) = {r ∈ R : rs = 0} is a maximal left ideal of R; S ≅ R/ann(s).
- $\operatorname{Rad}(R)$ annihilates simple modules: $\operatorname{Rad}(R)S = 0$.
- Every simple module is a module over R/Rad(R).
- Soc(A) is a module over $R/\operatorname{Rad}(R)$.
- Same idea for right modules; reverse sides.

Top-bottom duality

- ▶ *R* finite ring with 1; *A* finite left *R*-module.
- A/Rad(R)A is the "top quotient" of A; it is a sum of simple modules.
- $\operatorname{Soc}(\widehat{A}) = (\widehat{A} : \operatorname{Rad}(R)A) \cong (A/\operatorname{Rad}(R)A)^{\widehat{}}.$
- ▶ \supseteq : $(A/\operatorname{Rad}(R)A)^{\frown}$ is a sum of simple modules.
- \subseteq : because $Soc(\widehat{A}) Rad(R) = 0$.

Sketch of proof

- We already know $1 \leftrightarrow 2$.
- Fact: if $R = M_{k \times k}(\mathbb{F}_q)$, then $\widehat{R} \cong R$.
- Then general $(R/\operatorname{Rad}(R))^{\widehat{}} \cong R/\operatorname{Rad}(R)$.
- So $\operatorname{Soc}(\widehat{R}) \cong (R/\operatorname{Rad}(R))^{\widehat{}} \cong R/\operatorname{Rad}(R).$
- ▶ 1,2 ⇒ 3: If $\widehat{R} \cong R$, then Soc(R) \cong Soc(\widehat{R}) \cong R/Rad(R).

Construction

- $M_{k \times k}(\mathbb{F}_q)$ has a generating character: $\varrho(P) = \vartheta_q(\operatorname{Tr} P), P \in M_{k \times k}(\mathbb{F}_q).$
- ▶ Tr P is the matrix trace of P.
- If $q = p^e$ and $x \in \mathbb{F}_q$, then

$$\vartheta_q(x) = (x + x^p + \cdots x^{p^{e^{-1}}})/p \in \mathbb{Q}/\mathbb{Z}.$$

• ϑ_q is a generating character of \mathbb{F}_q .

Why does ϱ generate?

- Suppose $B \subseteq \ker \rho$ is a left ideal of R.
- Then $Soc(B) = B \cap Soc(R) \subseteq \ker \varrho \cap Soc(R)$.
- But *ρ* is a generating character of Soc(*R*), so Soc(*B*) = 0.
- Thus B = 0; ϱ is a left generating character of R.

Similar characterization for modules

Theorem

The following are equivalent:

- 1. \widehat{A} is a cyclic right R-module.
- 2. A injects into \widehat{R} : $A \hookrightarrow \widehat{R}$.
- 3. There exists $\varrho \in \widehat{A}$ such that ker ϱ contains no nonzero left *R*-submodule.
- 4. Soc(A) \subseteq A is a cyclic R-submodule.

More identifications

- R finite Frobenius ring with generating character ρ .
- Dot product on R^n : $y \cdot x = \sum_{i=1}^n y_i x_i$.
- Define $\psi: \mathbb{R}^n \to \widehat{\mathbb{R}}^n$, $x \mapsto \psi_x$:

$$\psi_x(y) = \varrho(y \cdot x), \quad y \in R^n.$$

• Then ψ is an isomorphism of left *R*-modules.

•
$$\psi_{rx}(y) = \varrho(y \cdot rx) = \varrho(yr \cdot x) = \psi_x(yr) = (r\psi_x)(y).$$

Character annihilator vs. dot product

• Recall:
$$\psi_x(y) = \varrho(y \cdot x), \quad y \in R^n.$$

- Additive subgroup $C \subseteq R^n$. Under ψ , $(\widehat{R}^n : C)$ corresponds to $r_{\varrho}(C) = \{x \in R^n : \varrho(C \cdot x) = 0\}$.
- Set $r(C) = \{x \in R^n : C \cdot x = 0\}.$
- $r(C) \subseteq r_{\varrho}(C)$ in general
- $r(C) = r(RC) = r_{\varrho}(RC) \subseteq r_{\varrho}(C)$ in general.
- r(C) = r_ℓ(C) when C is a left submodule, as C · x is a left ideal in ker ℓ.

Binary case

- Let q = 2, the binary case.
- For $x \in \mathbb{F}_2^n$, if $x \cdot x = 0$, then wt(x) is even. (This is also true for q = 3, but not in general.)
- If C ⊆ 𝔽ⁿ₂ is self-orthogonal, then every codeword in C has even weight.
- Extra: a binary self-orthogonal code in which every codeword has weight divisible by 4 is **doubly-even** (singly-even otherwise).

A binary example

► The codes generated by *G*₂, *G*₈ are singly-even, self-dual:

$$G_2 = [1 \ 1], \quad G_8 = egin{bmatrix} 1 & 1 & 0 & 0 & 0 & 0 & 0 \ 0 & 0 & 1 & 1 & 0 & 0 & 0 \ 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 \ 0 & 0 & 0 & 0 & 0 & 1 & 1 \end{bmatrix}$$

- hwe $_{G_2} = X^2 + Y^2$.
- hwe_{G₈} = $X^8 + 4X^6Y^2 + 6X^4Y^4 + 4X^2Y^6 + Y^8 = (X^2 + Y^2)^4$.

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Another binary example

• The code generated by E_8 is doubly-even, self-dual.

• hwe_{*E*₈} = $X^8 + 14X^4Y^4 + Y^8$.

Binary self-dual case

When the code C is self-dual, C appears on both sides of the MacWilliams identities:

$$\mathsf{hwe}_{C}(X, Y) = \frac{1}{|C|} \mathsf{hwe}_{C}(X + Y, X - Y).$$

Length is n = 2k. hwe_C(X, Y) is a homogeneous polynomial of degree n, so

$$hwe_{\mathcal{C}}(X,Y) = hwe_{\mathcal{C}}\left(\frac{X+Y}{\sqrt{2}},\frac{X-Y}{\sqrt{2}}\right)$$

Invariance properties

► The group GL(2, C) acts on C[X, Y] by linear substitution:

$$f(X, Y) \begin{bmatrix} a & b \\ c & d \end{bmatrix} = f(aX + cY, bX + dY).$$

For binary self-dual C, hwe_C is invariant under

$$M = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}.$$

More invariance properties

In addition, singly-even and doubly-even are invariant under, respectively (i = √−1):

$$W_s = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}, \quad W_d = \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix}$$

- Define two subgroups of GL(2, ℂ): G_s = ⟨M, W_s⟩ and G_d = ⟨M, W_d⟩.
- For singly-even C, hwe_C $\in \mathbb{C}[X, Y]^{\mathcal{G}_s}$.
- For doubly-even *C*, hwe_{*C*} $\in \mathbb{C}[X, Y]^{\mathcal{G}_d}$.

Examples

- Let S be a ring with anti-isomorphism ϵ .
- For any finite group G, the group ring R = S[G] has anti-isomorphism ε:

$$arepsilon(\sum_{g\in G} c_g g) = \sum_{g\in G} \epsilon(c_g) g^{-1}.$$

• Matrix ring $R = M_{k \times k}(S)$, using the transpose:

$$\varepsilon(P) = (\epsilon(P))^T, P \in R.$$

Apply ϵ to each entry of P.

Swapping sides

- ► An anti-isomorphism ε on R allows one to regard left modules as right modules, and vice versa.
- If M is a left R-module, define ε(M) to be same abelian group as M, but equipped with right scalar multiplication defined by

$$xr = \varepsilon(r)x, \quad x \in M, r \in R,$$

where $\varepsilon(r)x$ is the left scalar multiplication of the module M.

Similar definition for right module to left.

Interpret in terms of bi-additive form

- Use the additive form of characters:
 Â = Hom_ℤ(A, ℚ/ℤ).
- Define β : A × A → Q/Z by β(a, b) = ψ(b)(a), for a, b ∈ A. Extend additively to Aⁿ × Aⁿ. Then:
- β is bi-additive.
- $\beta(rx, y) = \beta(x, \varepsilon(r)y)$ for $x, y \in A^n$, $r \in R$.
- Impose one more property: there exists a unit e ∈ R such that β(x, y) = β(ey, x) for x, y ∈ Aⁿ.

Properties of C^{\perp}

• Recall
$$C^{\perp} = \psi^{-1}(\widehat{A}^n : C).$$

- In terms of β : $C^{\perp} = \{y \in A^n : \beta(C, y) = 0\}.$
- Even if $C \subseteq A^n$ is just an additive code, we have $|C| \cdot |C^{\perp}| = |A^n|$ and the MacWilliams identities.
- If C is a left linear code, then so is C^{\perp} .
- If C is a left linear code, then (C[⊥])[⊥] = C. This uses the β(x, y) = β(ey, x) condition.
- When C is a left linear code, we also have $C^{\perp} = \{x \in A^n : \beta(x, C) = 0\}.$

Example (c)

For k = 2, there are proper left ideals (a, b ∈ 𝔽₂):
C₁ = { [a 0 | b 0] }, C₂ = { [0 a | 0 b] }, C₃ = { [a a | b b] }.
Then C₁[⊥] = C₂, C₂[⊥] = C₁, and C₃[⊥] = C₃.