

Character-Theoretic Tools for Studying Linear Codes over Rings and Modules

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2. Additive Codes and Characters

- ▶ Definitions
- ▶ Properties
- ▶ Fourier transform
- ▶ MacWilliams identities
- ▶ **Exercises** throughout

Additive codes

- ▶ Let A be a finite abelian group (additive notation); A will be a module later.
- ▶ An **additive code** of length n over A is an additive subgroup $C \subseteq A^n$.
- ▶ The **Hamming weight** on A , $\text{wt} : A \rightarrow \mathbb{C}$, is

$$\text{wt}(a) = \begin{cases} 0, & a = 0, \\ 1, & a \neq 0. \end{cases}$$

- ▶ Extend to A^n by $\text{wt}(a_1, \dots, a_n) = \sum \text{wt}(a_i)$.

Hamming weight enumerator

- ▶ For an additive code $C \subseteq A^n$, define the **Hamming weight enumerator** of C by

$$\text{hwe}_C(X, Y) = \sum_{x \in C} X^{n-\text{wt}(x)} Y^{\text{wt}(x)}.$$

- ▶ $\text{hwe}_C(X, Y) = \sum_{i=0}^n A_i X^{n-i} Y^i$, where A_i is the number of codewords in C of Hamming weight i .

How to form a dual code?

- ▶ We would like to form a dual code, but there is no dot product immediately available.
- ▶ Form a dual code abstractly!

Characters

- ▶ A **character** of A is a group homomorphism

$$\pi : A \rightarrow \mathbb{C}^\times,$$

where \mathbb{C}^\times is the multiplicative group of nonzero complex numbers: $\pi(a + b) = \pi(a)\pi(b)$, $a, b \in A$.

- ▶ *Representation theory: π is the character of a 1-dimensional complex representation of A . Because A is abelian, every irreducible complex representation of A is 1-dimensional.*

Character group

- ▶ The set \widehat{A} of all characters of A is a multiplicative abelian group under pointwise multiplication.

$$(\pi\psi)(a) = \pi(a)\psi(a), \quad a \in A, \quad \pi, \psi \in \widehat{A}.$$

- ▶ **Exercise:** every character of $\mathbb{Z}/k\mathbb{Z}$ has the form $\rho_b(a) = \exp(2\pi iab/k)$, $a \in \mathbb{Z}/k\mathbb{Z}$, for some $b \in \mathbb{Z}/k\mathbb{Z}$. [What is $\rho(1)$?]
- ▶ Thus, $(\mathbb{Z}/k\mathbb{Z})^\wedge \cong \mathbb{Z}/k\mathbb{Z}$, via $\rho_b \longleftrightarrow b$.

Additive form of character group

- ▶ Original, multiplicative form: $\widehat{A} = \text{Hom}_{\mathbb{Z}}(A, \mathbb{C}^{\times})$.
- ▶ Additive version: $\widehat{A} \cong \text{Hom}_{\mathbb{Z}}(A, \mathbb{Q}/\mathbb{Z})$.
- ▶ $\varrho \in \text{Hom}_{\mathbb{Z}}(A, \mathbb{Q}/\mathbb{Z})$ corresponds to $\rho \in \text{Hom}_{\mathbb{Z}}(A, \mathbb{C}^{\times})$ by $\rho(a) = \exp(2\pi i \varrho(a))$.
- ▶ $\rho(a + b) = \rho(a)\rho(b)$, while $\varrho(a + b) = \varrho(a) + \varrho(b)$.

Duality functor

- ▶ Pontryagin duality: $A \mapsto \widehat{A}$
- ▶ Exact contravariant functor:

$$0 \rightarrow A_1 \rightarrow A_2 \rightarrow A_3 \rightarrow 0$$

induces

$$0 \rightarrow \widehat{A}_3 \rightarrow \widehat{A}_2 \rightarrow \widehat{A}_1 \rightarrow 0.$$

- ▶ $\widehat{\widehat{A}} \cong A$, but not naturally. (*Uses fundamental theorem of finitely generated abelian groups.*)
- ▶ $\widehat{\widehat{\widehat{A}}} \cong A$, naturally: $a \mapsto (\pi \mapsto \pi(a))$.
- ▶ $(A \times B)^\wedge \cong \widehat{A} \times \widehat{B}$.

Annihilators

- ▶ Let $B \subseteq A$ be any subgroup.
- ▶ Define the **annihilator** $(\hat{A} : B)$:

$$(\hat{A} : B) = \{\rho \in \hat{A} : \rho(B) = 1\} = \{\varrho \in \hat{A} : \varrho(B) = 0\}.$$

- ▶ $(\hat{A} : B) \cong (A/B)^\wedge$.
- ▶ $|B| \cdot |(\hat{A} : B)| = |A|$.
- ▶ Double annihilator: $(A : (\hat{A} : B)) = B$.

Application to additive codes

- ▶ Let A be a finite abelian group, and let $C \subseteq A^n$ be an additive code.
- ▶ View $C \subseteq A^n$ as an example of “ $B \subseteq A^n$ ”.
- ▶ The **dual code** of $C \subseteq A^n$ is the annihilator $(\widehat{A}^n : C) \subseteq \widehat{A}^n$.

Good duality properties

- ▶ Given an additive code $C \subseteq A^n$.
- ▶ Dual $(\widehat{A}^n : C) \subseteq \widehat{A}^n$ is an additive code over \widehat{A} .
- ▶ Double annihilator: $(A^n : (\widehat{A}^n : C)) = C$.
- ▶ Size: $|C| \cdot |(\widehat{A}^n : C)| = |A^n|$.
- ▶ The MacWilliams identities. (Coming next.)

Two weight enumerators

- ▶ The Hamming weight enumerator of C is

$$\text{hwe}_C(X, Y) = \sum_{x \in C} X^{n-\text{wt}(x)} Y^{\text{wt}(x)}.$$

- ▶ The **complete weight enumerator** of C is a homogeneous polynomial in $\mathbb{C}[Z_a : a \in A]$:

$$\text{cwe}_C((Z_a)) = \sum_{x \in C} \prod_{i=1}^n z_{x_i}.$$

MacWilliams Identities

- ▶ The MacWilliams identities express the Hamming or complete weight enumerators of C in terms of those of its dual code $(A^n : C)$.
- ▶ The expression involves a linear change of variables.
- ▶ The Hamming case, with $C^\perp = (\hat{A}^n : C)$:

$$\text{hwe}_C(X, Y) = \frac{1}{|C^\perp|} \text{hwe}_{C^\perp}(X + (|A| - 1)Y, X - Y).$$

- ▶ Proof involves the Fourier transform.

Summation formulas

- ▶ Need multiplicative form of characters.
- ▶ For $\pi \in \widehat{A}$,

$$\sum_{a \in A} \pi(a) = \begin{cases} |A|, & \pi = 1, \\ 0, & \pi \neq 1. \end{cases}$$

- ▶ For $a \in A$,

$$\sum_{\pi \in \widehat{A}} \pi(a) = \begin{cases} |A|, & a = 0, \\ 0, & a \neq 0. \end{cases}$$

Fourier transform

- ▶ Given a function $f : A \rightarrow V$, V a complex vector space. Define its **Fourier transform** $\hat{f} : \hat{A} \rightarrow V$ by

$$\hat{f}(\pi) = \sum_{a \in A} \pi(a) f(a), \quad \pi \in \hat{A}.$$

- ▶ $\hat{\cdot} : F(A, V) \rightarrow F(\hat{A}, V)$.
- ▶ Invert:

$$f(a) = \frac{1}{|A|} \sum_{\pi \in \hat{A}} \pi(-a) \hat{f}(\pi), \quad a \in A.$$

Poisson summation formula

Let B be any subgroup of A , and let $f : A \rightarrow V$. Then for any $a \in A$,

$$\sum_{b \in B} f(a + b) = \frac{1}{|(\widehat{A} : B)|} \sum_{\pi \in (\widehat{A} : B)} \pi(-a) \widehat{f}(\pi).$$

If $a = 0$, then

$$\sum_{b \in B} f(b) = \frac{1}{|(\widehat{A} : B)|} \sum_{\pi \in (\widehat{A} : B)} \widehat{f}(\pi).$$

A Fourier transform example

- ▶ Suppose V is a complex algebra.
- ▶ Suppose $f : A^n \rightarrow V$ has the form

$$f(a_1, \dots, a_n) = \prod_{i=1}^n f_i(a_i),$$

where $f_i : A \rightarrow V$.

- ▶ Then

$$\hat{f}(\pi_1, \dots, \pi_n) = \prod_{i=1}^n \hat{f}_i(\pi_i).$$

Complete weight enumerator

- ▶ $V = \mathbb{C}[Z_a : a \in A]$, a complex algebra.
- ▶ $f : A^n \rightarrow V$,

$$f(a_1, \dots, a_n) = \prod_{i=1}^n Z_{a_i}.$$

- ▶ Then

$$\hat{f}(\pi_1, \dots, \pi_n) = \prod_{i=1}^n \left(\sum_{a_i \in A} \pi_i(a_i) Z_{a_i} \right)$$

MacWilliams identities from Poisson summation formula

- ▶ Poisson:

$$\sum_{b \in B} f(b) = \frac{1}{|(\widehat{A} : B)|} \sum_{\pi \in (\widehat{A} : B)} \widehat{f}(\pi).$$

- ▶ Replace A by A^n , B by additive code C , $(\widehat{A} : B)$ by dual code $(\widehat{A}^n : C)$.

MacWilliams identities: complete weight enumerator

- ▶ $Z = (Z_a)_{a \in A}$; $f(a_1, \dots, a_n) = \prod_{i=1}^n Z_{a_i}$.
- ▶ Complete weight enumerator:

$$\text{cwe}_C(Z) = \sum_{x \in C} f(x) = \sum_{a \in C} \prod_{i=1}^n Z_{a_i}.$$

- ▶ MacWilliams identities:

$$\text{cwe}_C(Z) = \frac{1}{|(\widehat{A}^n : C)|} \text{cwe}_{(\widehat{A}^n : C)}\left(\sum_{a \in A} \pi(a) Z_a\right).$$

Specialize to Hamming weight enumerator

- ▶ Recall $\text{hwe}_C(X, Y) = \sum_{x \in C} X^{n-\text{wt}(x)} Y^{\text{wt}(x)}$.
- ▶ Specialize $\mathbb{C}[Z_a : a \in A] \rightarrow \mathbb{C}[X, Y]$, $Z_0 \mapsto X$, $Z_a \mapsto Y$ for $a \neq 0$. Then $\text{cwe}_C(Z) \mapsto \text{hwe}_C(X, Y)$.
- ▶ What happens to $\text{cwe}_{(\hat{A}^n: C)}(\sum_{a \in A} \pi(a) Z_a)$ on the right side?

Specialization

$$\sum_{a \in A} \pi(a) Z_a = \pi(0) Z_0 + \sum_{a \neq 0} \pi(a) Z_a$$

$$\mapsto X + \left(\sum_{a \neq 0} \pi(a) \right) Y$$

$$= \begin{cases} X + (|A| - 1)Y, & \text{if } \pi = 1, \\ X - Y, & \text{if } \pi \neq 1. \end{cases}$$

MacWilliams identities: Hamming weight enumerator

$$\text{cwe}_C(Z) = \frac{1}{|(\widehat{A}^n : C)|} \text{cwe}_{(\widehat{A}^n : C)}\left(\sum_{a \in A} \pi(a) Z_a\right)$$

specializes to

$$\text{hwe}_C(X, Y) = \frac{1}{|C^\perp|} \text{hwe}_{C^\perp}(X + (|A| - 1)Y, X - Y),$$

where $C^\perp = (\widehat{A}^n : C)$.

Next steps

- ▶ What happens when A is a left module over a finite ring R and $C \subseteq A^n$ is a linear code?
- ▶ Is the dual code $(\widehat{A}^n : C)$ linear?
- ▶ What duality properties hold?
- ▶ If $A = R$, can $(\widehat{R}^n : C)$ be expressed in terms of the dot product on R^n ?