

Random Graphs III

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Outline of Lecture III

1. **Subgraph containment with adversary:**

- ▷ Existence of mono_χ subgraphs in **coloured random graphs**; properties of the form $G(n, p) \rightarrow_r (H)_r^e$ (omit “e”; we are done with vertex colourings)
- ▷ Existence of monochromatic subgraphs in ‘dense’ subgraphs of random graphs; statements of the form $G(n, p) \rightarrow_\eta H$

2. **Regularity method for sparse graphs:** some of the (partial) successes

3. **Literature:** shall discuss some of the literature

One last thing about the vertex case

Exercise 1⁺⁺: determine $p_0 = p_0(n)$ such that if $p \gg p_0$, then a.e. $G(n, p)$ is such that

$$G(n, p) \rightarrow (K^3, K^4)^v \quad (1)$$

Have you obtained the threshold?

The results for K^3

Theorem 1. *There is a large enough constant C such that if $p \geq C/\sqrt{n}$, then a.e. $G(n, p)$ satisfies*

$$G(n, p) \rightarrow (K^3)_2. \quad (2)$$

Theorem 2. *For any $\eta > 0$, there C such that if $p \geq C/\sqrt{n}$ then a.e. $G(n, p)$ satisfies*

$$G(n, p) \rightarrow_{1/2+\eta} K^3. \quad (3)$$

The case $p = 1$

Theorem 3 (Goodman). *The number of monochromatic triangles in any 2-colouring of K^n is $\geq \binom{n}{3} - \frac{1}{8}n(n-1)^2 = (\frac{1}{4} + o(1))\binom{n}{3}$.*

Proof. Count the number N of 2-coloured cherries. Say have $b(x)$ blue edges and $r(x)$ red edges at vertex x . Have $N = \sum_x b(x)r(x) \leq n(n-1)^2/4$ (use $b(x) + r(x) = n-1$). The number of non-monochromatic triangles is $N/2$. \square

Clearly, there are 2-colourings with $\leq \frac{1}{4}\binom{n}{3}$ monochromatic triangles!

The case $p = 1$

Theorem 4. *Goodman implies Mantel: $ex(n, K^3) \leq n^2/4$.*

Proof. Suppose $K^3 \not\subset G^n$ and $e(G^n) > n^2/4$, so that $e(G^n) \geq n^2/4 + 3/4$. Count the number N of pairs (e, T) where $e \in E(G^n)$, $T \in \binom{V(G^n)}{3}$ and $e \subset T$. Say there are t_i triples containing i edges of G . Then, by Theorem 3,

$$\begin{aligned} (n^2/4 + 3/4)(n - 2) &\leq e(G^n)(n - 2) \\ &= N \leq t_1 + 2t_2 \leq 2(t_1 + t_2) \leq n(n - 1)^2/4, \end{aligned}$$

which is a contradiction for $n \geq 4$. □

Theorems 1 and 2 from Goodman's counting

Theorem 5. *For every $\varepsilon > 0$ there is C such that if $p \geq C/\sqrt{n}$, then a.e. $G(n, p)$ is such that any 2-colouring of its edges contains at least $(1/4 - \varepsilon)p^3 \binom{n}{3}$ monochromatic K^3 .*

Proof. Exercise 2⁺. □

Exercise 3⁺: derive Theorem 2 from Theorem 5.

Exercise 4⁺⁺: show that the threshold for Theorem 1 is indeed $1/\sqrt{n}$.

Exercise 5: show that the threshold for Theorem 2 is indeed $1/\sqrt{n}$.

The Rödl–Ruciński theorem

Definition 6 (2-density and $m_2(H)$). *The 2-density $d_2(H)$ of a graph H with $|V(H)| > 2$ is*

$$\frac{|E(H)| - 1}{|V(H)| - 2}. \quad (4)$$

For $H = K^1$ and $2K^1$ let $d_2(H) = 0$; set $d_2(K^2) = 1/2$. Let

$$m_2(H) = \max\{d_2(J) : J \subset H, |V(J)| > 0\}. \quad (5)$$

Exercise 6: consider, say, $H = K^h$. Show that if $p \ll n^{-1/m_2(H)}$, then a.s. $\#\{H \hookrightarrow G(n, p)\} \ll e(G(n, p))$. On the other hand, if $p \gg n^{-1/m_2(H)}$, then a.s. $\#\{H \hookrightarrow G(n, p)\} \gg e(G(n, p))$.

The Rödl–Ruciński theorem

Theorem 7. *Let H be a graph containing at least a cycle and let $r \geq 2$ be an integer. Then there exist constants c and C such that*

$$\lim_{n \rightarrow \infty} \mathbb{P}(G(n, p) \rightarrow (H)_r) = \begin{cases} 0 & \text{if } p \leq cn^{-1/m_2(H)} \\ 1 & \text{if } p \geq Cn^{-1/m_2(H)}. \end{cases} \quad (6)$$

▷ In particular, $p_0 = p_0(n) = n^{-1/m_2(H)}$ is a threshold for the property $G(n, p) \rightarrow (H)_r$.

The Rödl–Ruciński theorem

Exercise 7: show that there exists a graph G with the property that $G \rightarrow (K^h)_r$ but $K^{h+1} \not\subseteq G$.

Exercise 8: Given a graph G , let $\mathcal{H}_3(G)$ be the 3-uniform hypergraph whose hypervertices are the edges of G and the hyperedges are the edge sets of the triangles in G . Show that, for any integers ℓ and r , there is a graph G satisfying $G \rightarrow (K^3)_r$ such that $\mathcal{H}_3(G)$ has girth $\geq \ell$.

Turán type results for subgraphs of random graphs

Generalized Turán number:

$$\text{ex}(G, H) = \max \{ |E(G')| : H \not\subset G' \subset G \}. \quad (7)$$

- ▷ $\text{ex}(n, H) = \text{ex}(K^n, H)$
- ▷ $\text{ex}(Q^d, C^4) = ?$
- ▷ $\text{ex}(G, K^h) = ?$ for (n, d, λ) -graphs G

Turán type results for subgraphs of random graphs

Exercise 9: show that, for any G and H , we have

$$e(G, H) \geq \left(1 - \frac{1}{\chi(H) - 1}\right) e(G). \quad (8)$$

▷ Interested in knowing when this is sharp for $G = G(n, p)$ (up to $o(e(G))$)

Turán type results for subgraphs of random graphs

Theorem 8. *Let H be a graph with degeneracy d and suppose $np^d \gg 1$. Then*

$$\text{ex}(G(n, p), H) = \left(1 - \frac{1}{\chi(H) - 1} + o(1)\right) p \binom{n}{2}. \quad (9)$$

Conjecture 9. *For any graph H , if $np^{m_2(H)} \rightarrow \infty$ then (9) holds almost surely.*

Example: if $H = K^k$, have $m_2(H) = (k + 1)/2$, but Theorem 8 supposes $np^{1/(k-1)} \gg 1$.

Cycles and small cliques

Theorem 10. *Conjecture 9 holds for cycles.*

Theorem 11. *Conjecture 9 holds for K^4 , K^5 , and K^6 .*

A sharp result for even cycles

Theorem 12. *Let $k \geq 2$ be an integer and let $p = p(n) = \alpha n^{-1+1/(2k-1)}$ be such that*

$$2 \leq \alpha \leq n^{1/(2k-1)^2}. \quad (10)$$

Then

$$\text{ex}(G(n, p), C^{2k}) \asymp \frac{(\log \alpha)^{1/(2k-1)}}{\alpha} e(G(n, p)). \quad (11)$$

Szemerédi's regularity lemma

1. Tool for identifying the quasirandom structure of deterministic graphs
2. Works very well for large, **dense** graphs: **n -vertex graphs with $\geq cn^2$ edges, $n \rightarrow \infty$**
3. Variant for sparse graphs exists (sparse = with $o(n^2)$ edges)
4. Much harder to use

ε -regularity

Definition 13 (ε -regular pair). $G = (V, E)$ a graph; $U, W \subset V$ non-empty and disjoint. Say (U, W) is ε -regular (in G) if

▷ for all $U' \subset U, W' \subset W$ with $|U'| \geq \varepsilon|U|$ and $|W'| \geq \varepsilon|W|$, we have

$$\left| \frac{|E(U', W')|}{|U'||W'|} - \frac{|E(U, W)|}{|U||W|} \right| \leq \varepsilon. \quad (12)$$

ε -regularity

The pair (U, W) is ε -regular if

$$\left| \frac{|E(U', W')|}{|U'| |W'|} - \frac{|E(U, W)|}{|U| |W|} \right| \leq \varepsilon. \quad (13)$$

Equivalently,

$$|E(U', W')| = |U'| |W'| \left(\frac{|E(U, W)|}{|U| |W|} \pm \varepsilon \right) \quad (14)$$

Clearly, **not meaningful** if

$$\frac{|E(U, W)|}{|U| |W|} \rightarrow 0 \quad (15)$$

and ε is fixed. (We think of $G = (V, E)$ with $n = |V| \rightarrow \infty$.)

ε -regularity; multiplicative error version

Replace

$$|E(U', W')| = |U'| |W'| \left(\frac{|E(U, W)|}{|U| |W|} \pm \varepsilon \right) \quad (16)$$

by

$$|E(U', W')| = (1 \pm \varepsilon) |E(U, W)| \frac{|U'| |W'|}{|U| |W|} \quad (17)$$

Altered condition becomes

▷ for all $U' \subset U$, $W' \subset W$ with $|U'| \geq \varepsilon |U|$ and $|W'| \geq \varepsilon |W|$, we have

$$\left| |E(U', W')| - |E(U, W)| \frac{|U'| |W'|}{|U| |W|} \right| \leq \varepsilon |E(U, W)| \frac{|U'| |W'|}{|U| |W|}. \quad (18)$$

OK even if $|E(U, W)| / |U| |W| \rightarrow 0$.

Szemerédi's regularity lemma, sparse version

Any graph with no 'dense patches' admits a Szemerédi partition with the new notion of ε -regularity.

Definition 14 ((η, b) -bounded). Say $G = (V, E)$ is (η, b) -bounded if for all $U \subset V$ with $|U| \geq \eta|V|$, we have

$$\#\{\text{edges within } U\} \leq b|E| \binom{|U|}{2} \binom{|V|}{2}^{-1}. \quad (19)$$

[▷ Something like one-sided (p, η) -uniformity]

Szemerédi's regularity lemma, sparse version

Theorem 15 (The regularity lemma). *For any $\varepsilon > 0$, $t_0 \geq 1$, and b , there exist $\eta > 0$ and T_0 such that any (η, b) -bounded graph G admits a partition $V = V_1 \cup \dots \cup V_t$ such that*

1. $|V_1| \leq \dots \leq |V_t| \leq |V_1| + 1$
2. $t_0 \leq t \leq T_0$
3. *at least $(1 - \varepsilon) \binom{t}{2}$ pairs (V_i, V_j) ($i < j$) are ε -regular.*

Proof. Just follow Szemerédi's original proof. □

A counting lemma (simplest version)

Setup. $G = (V_1, V_2, V_3; E)$ tripartite with

1. $|V_i| = m$ for all i
2. (V_i, V_j) ε -regular for all $i < j$
3. $|E(V_i, V_j)| = \rho m^2$ for all $i < j$

Notation: $G = G_3^{(\varepsilon)}(m, \rho)$ [G is an ε -regular triple with density ρ]

▷ Wish to embed K^3 with $V(K^3) = \{x_1, x_2, x_3\}$ such that x_i is placed in V_i .

A counting lemma (simplest version)

Just like random:

Lemma 16 (Counting Lemma; Embedding lemma). $\forall \rho > 0, \delta > 0 \exists \varepsilon > 0,$
 m_0 : if $m \geq m_0$, then

$$\left| \#\{K^3 \hookrightarrow G_3^{(\varepsilon)}(m, \rho)\} - \rho^3 m^3 \right| \leq \delta m^3. \quad (20)$$

Tough: counting Lemma is false if $\rho \rightarrow 0$

Fact 17. $\forall \varepsilon > 0 \exists \rho > 0, m_0 \forall m \geq m_0 \exists G_3^{(\varepsilon)}(m, \rho)$ with

$$K^3 \not\subset G_3^{(\varepsilon)}(m, \rho). \quad (21)$$

[cf. Lemma 16]

Change of focus: from *counting* $K^3 \subset G_3^{(\varepsilon)}(m, \rho)$ to *existence of* $K^3 \subset G_3^{(\varepsilon)}(m, \rho)$

Key observation

Counterexamples to the embedding lemma in the sparse setting do exist (Fact 17), but

are extremely rare.

An asymptotic enumeration lemma

Lemma 18. $\forall \beta > 0 \exists \varepsilon > 0, C > 0, m_0$: if $T = \rho m^2 \geq Cm^{3/2}$, then

$$\#\{G_3^{(\varepsilon)}(m, \rho) \not\cong K^3\} \leq \beta T \binom{m^2}{T}^3. \quad (22)$$

Observe that $\rho \geq C/\sqrt{m} \rightarrow 0$.

Consequence for random graphs

Easy expectation calculations imply

▷ if $p \gg 1/\sqrt{n}$, then almost every $G(n, p)$ is such that

$$\left(K^3\text{-free } G_3^{(\varepsilon)}(m, \rho) \right) \not\subset G(n, p), \quad (23)$$

if (*) $mp \gg \log n$ and $\rho \geq \alpha p$ for some fixed α .

Conclusion. Recovered an ‘embedding lemma’ in the sparse setting, *for subgraphs of random graphs*.

Corollary 19 (EL for subgraphs of r.gs). *If $p \gg 1/\sqrt{n}$ and (*) holds, then almost every $G(n, p)$ is such that if $G_3^{(\varepsilon)}(m, \rho) \subset G(n, p)$, then*

$$\exists \iota: K^3 \hookrightarrow G_3^{(\varepsilon)}(m, \rho) \subset G(n, p). \quad (24)$$

The easy calculation

Recall $T = \rho m^2$, $\rho \geq \alpha p$, and $mp \gg \log n$. Therefore

$$\begin{aligned}
 \mathbb{E}(\#\{K^3\text{-free } G_3^{(\varepsilon)}(m, \rho) \hookrightarrow G(n, p)\}) &\leq n^{3m} o(1)^T \binom{m^2}{T}^3 p^{3T} \\
 &\leq n^{3m} \left(o(1) \frac{em^2}{T}\right)^{3T} p^{3T} \leq n^{3m} \left(o(1) \frac{em^2 p}{T}\right)^{3T} \\
 &= n^{3m} \left(o(1) \frac{em^2 p}{\rho m^2}\right)^{3T} \leq n^{3m} \left(o(1) \frac{e}{\alpha}\right)^{3T} \\
 &\leq e^{3m \log n} \left(o(1) \frac{e}{\alpha}\right)^{3\alpha m \log n} = o(1).
 \end{aligned}$$

Superexponential bounds

Suppose we wish to prove a statement about **all subgraphs** of $G(n, p)$.

- Too many such subgraphs: about $2^p \binom{n}{2}$
- $G(n, p)$ has no edges with probability $(1 - p)^{\binom{n}{2}} \geq \exp\{-2pn^2\}$, if, say, $p \leq 1/2$.
- Concentration inequalities won't do ($2^p \binom{n}{2}$ vs e^{-2pn^2}).
- Bounds of the form

$$o(1)^T \binom{m}{T} \tag{25}$$

for the cardinality of a family of **'undesirable subgraphs'** $\mathcal{U}(m, T)$ do the job. Use of such bounds goes back to Füredi (1994).

An embedding lemma for K^3 (sparse setting)

Scheme:

1. Proved an asymptotic enumeration lemma:

$$\#\{G_3^{(\varepsilon)}(m, \rho) \not\supset K^3\} = o(1) \binom{m^2}{T}^3. \quad (26)$$

$$[T = \rho m^2]$$

2. Observed that this implies a.e. $G(n, \omega/\sqrt{n})$ contains no K^3 -free $G_3^{(\varepsilon)}(m, \rho)$ [any $\omega = \omega(n) \rightarrow \infty$ as $n \rightarrow \infty$].
3. Obtained a K^3 -embedding lemma for subgraphs of $G(n, p)$, even when $p = \omega/\sqrt{n}$: for all $G_3^{(\varepsilon)}(m, \rho) \subset G(n, p)$, have $K^3 \subset G_3^{(\varepsilon)}(m, \rho)$.

General graphs H ?

Let us consider the case $H = K^k$.

Conjecture 20. $\forall k \geq 4, \beta > 0 \exists \varepsilon > 0, C > 0, m_0$: if $T = \rho m^2 \geq Cm^{2-2/(k+1)}$, then

$$\#\{G_k^{(\varepsilon)}(m, \rho) \not\supseteq K^k\} \leq \beta T \binom{m^2}{T}^{(k)} . \quad (27)$$

▷ Known for $k = 3, k = 4$, and $k = 5$

▷ For general H , the conjecture involves $m_2(H)$

General graphs H ?

Suppose $|V(H)| > 2$. Then let

$$d_2(H) = \frac{|E(H)| - 1}{|V(H)| - 2}. \quad (28)$$

For $H = K^1$ and $2K^1$ let $d_2(H) = 0$; set $d_2(K^2) = 1/2$. Finally, let

$$m_2(H) = \max\{d_2(J) : J \subset H\}. \quad (29)$$

Conjecture 21. $\forall H, \beta > 0 \exists \varepsilon > 0, C > 0, m_0$: if $T = \rho m^2 \geq C m^{2-1/m_2(H)}$, then

$$\#\{G_H^{(\varepsilon)}(m, \rho) \not\supset H\} \leq \beta^T \binom{m^2}{T}^{e(H)}. \quad (30)$$

▷ Known for cycles

The bipartite case

Conjecture 22. *Suppose H is a bipartite graph. For any $\beta > 0$, there is C such that for any $M = M(n)$ such that $M \geq Cn^{2-1/m_2(H)}$, we have*

$$\#\{G(n, M) \not\supseteq H\} \leq \beta^M \binom{\binom{n}{2}}{M} \quad (31)$$

for all large enough n .

In short: $\#\{G(n, M) \not\supseteq H\} = o(1)^M \binom{\binom{n}{2}}{M}$

- Known for even cycles [**Exercise 10⁺⁺**: still **++**, even after regularity. What is simpler to show is that $\text{ex}(G(n, p), C^{2k}) = o(p \binom{\binom{n}{2}}{2k})$ if $p \gg n^{-1+1/(2k-1)}$.]

Consequences of the conjecture

1. [The Rödl–Ruciński theorem](#) on threshold for Ramsey properties of random graphs (1-statement) and the [Turán counterpart](#), with the best possible threshold
2. [Łuczak \(2000\)](#): structural and enumerative consequences for H -free graphs on n vertices and M edges

Consequences of the conjecture

Under the hypothesis that Conjecture 21 holds for a graph H :

Theorem 23. *Conjecture 22 holds: if H is bipartite, then for every $\beta > 0$ there exists $C = C(\beta, H)$ and n_0 such that for $n \geq n_0$ and $M \geq Cn^{2-1/m_2(H)}$ we have*

$$\#\{G(n, M) \not\supseteq H\} \leq \beta^M \binom{\binom{n}{2}}{M}. \quad (32)$$

Consequences of the conjecture

Under the hypothesis that Conjecture 21 holds for a graph H :

Theorem 24. *Suppose $\chi(H) = h \geq 3$. Then for every $\delta > 0$ there exists $C = C(\delta, H)$ such that, almost surely, a graph chosen uniformly at random from the family of all H -free labelled graphs on n vertices and $M \geq Cn^{2-1/m_2(H)}$ edges can be made $(h-1)$ -partite by removing $\leq \delta M$ edges.*

Consequences of the conjecture

Under the hypothesis that **Conjecture 21** holds for a graph H :

Theorem 25. *Suppose $\chi(H) = h \geq 3$. Then for every $\varepsilon > 0$ there exist $C = C(\varepsilon, H)$ and $n_0 = n_0(\varepsilon, H)$ such that, for $n \geq n_0$ and $Cn^{2-1/d_2(H)} \leq M \leq n^2/C$, we have*

$$\left(\frac{h-2}{h-1} - \varepsilon\right)^M \leq \mathbb{P}(G(n, M) \not\supseteq H) \leq \left(\frac{h-2}{h-1} + \varepsilon\right)^M. \quad (33)$$

The hereditary nature of regularity

Setup. $B = (U, W; E)$ an ε -regular bipartite graph with $|U| = |W| = m$ and $|E| = \rho m^2$, $\rho > 0$ constant, and an integer d . Sample $N \subset U$ and $N' \subset W$ with $|N| = |N'| = d$ uniformly at random.

Theorem 26. For any $\beta > 0$, $\rho > 0$, and $\varepsilon' > 0$, if $\varepsilon \leq \varepsilon_0(\beta, \rho, \varepsilon')$, $d \geq d_0(\beta, \rho, \varepsilon')$, and $m \geq m_0(\beta, \rho, \varepsilon')$, then

$$\mathbb{P}((N, N') \text{ bad}) \leq \beta^d, \quad (34)$$

where (N, N') is bad if $||E(N, N')|d^{-2} - \rho| > \varepsilon'$ or else (N, N') is not ε' -regular.

The hereditary nature of regularity

Exercise 11⁺⁺: for any k and $\delta > 0$, there is C such that the following holds. If $\chi(G - F) \geq k$ for any $F \subset E(G)$ with $|F| \leq \delta n^2$, then there is $H \subset G$ with $\chi(H) \geq k$ and $|V(H)| \leq C$. Can you guarantee many such 'witnesses' H ?

Local characterization for regularity

Setup. $B = (U, W; E)$, a bipartite graph with $|U| = |W| = m$. Consider the properties

(PC) for some constant p , have $m^{-1} \sum_{u \in U} |\deg(u) - pm| = o(m)$ and

$$\frac{1}{m^2} \sum_{u, u' \in U} |\deg(u, u') - p^2 m| = o(m). \quad (35)$$

(R) (U, W) is $o(1)$ -regular (classical sense).

Theorem 27. (PC) and (R) are equivalent.

A proof of Theorem 26

Let a graph $F = (U, E)$ with $|U| = m$ and $|E| \leq \eta \binom{m}{2}$ be given. Suppose we select a d -set N uniformly at random from U . We are then interested in giving an upper bound for $e(F[N])$, the number of edges that the set N will induce in F .

Lemma 28. *For every α and $\beta > 0$, there exist $\eta_0 = \eta_0(\alpha, \beta) > 0$ such that, whenever $0 < \eta \leq \eta_0$, we have*

$$\mathbb{P} \left(e(F[N]) \geq \alpha \binom{d}{2} \right) \leq \beta^d. \quad (36)$$

Proof. Exercise 12⁺. □

Exercise 13⁺⁺: use Lemma 28 and Theorem 27 to prove Theorem 26.

The hereditary nature of sparse regularity

Definition 29 ((ε, p) -lower-regularity). *Suppose $0 < \varepsilon < 1$ and $0 < p \leq 1$. A bipartite graph $B = (U, W; E)$ is (ε, p) -lower-regular if for all $U' \subset U$ and $W' \subset W$ with $|U'| \geq \varepsilon|U|$ and $|W'| \geq \varepsilon|W|$, we have*

$$\frac{e(U', W')}{|U'| |W'|} \geq (1 - \varepsilon)p. \quad (37)$$

The hereditary nature of sparse regularity

Setup. $B = (U, W; E)$ an (ε, p) -lower-regular bipartite graph with $|U| = |W| = m$ and $|E| = pm^2$ and integer d . Sample $N \subset U$ and $N' \subset W$ with $|N| = |N'| = d$ uniformly at random.

Theorem 30. *For all $0 < \beta, \varepsilon' < 1$, there exist $\varepsilon_0 = \varepsilon_0(\beta, \varepsilon') > 0$ and $C = C(\varepsilon')$ such that, for any $0 < \varepsilon \leq \varepsilon_0$ and $0 < p < 1$, the following holds. Let $G = (U, W; E)$ be an (ε, p) -lower-regular bipartite graph and suppose $d \geq Cp^{-1}$. Then*

$$\mathbb{P}((N, N') \text{ bad}) \leq \beta^d, \quad (38)$$

where (N, N') is bad if (N, N') is not (ε', p) -lower-regular.

The hereditary nature of sparse regularity

Setup. $B = (U, W; E)$ an (ε, p) -lower-regular bipartite graph with $|U| = |W| = m$, $|E| = pm^2$, and $p \geq \alpha q$. Also, suppose we have two other bipartite graphs $A = (U', U; E_A)$ and $C = (W, W'; E_C)$, also (ε, p) -lower-regular.

Corollary 31 (Quite imprecise...). *Suppose $A \cup B \cup C \subset G(n, q)$ and $u' \in U'$ and $w' \in W'$ are 'typical' vertices. If $pm \gg 1/p$ then $B[\Gamma_A(u'), \Gamma_C(w')]$ is $(f(\varepsilon), p)$ -lower-regular ($f(\varepsilon) \rightarrow 0$ as $\varepsilon \rightarrow 0$).*

▷ Corollary above may be used in inductive embedding schemes.

Local characterization for sparse regularity

Very imprecisely: a similar statement to Theorem 27 may be proved for subgraphs of random graphs.

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Yoshi and Yoshi, admiring your patience for being still in the
room

