

SOME INEQUALITIES WITH PROBABILITIES

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1. THE INEQUALITIES

Let C denote the set of all maps $\alpha : \{1, 2, 3\} \rightarrow \{0, 1\}$. We think of 0, 1 as *colors* (say, 0 for green and 1 for red) and of an element $\alpha \in C$ as a triple of colors. Given $\alpha, \beta \in C$, we set:

$$n(\alpha, \beta) = |\{(i, j) \in \{1, 2, 3\}^2 : i \neq j \text{ and } \alpha(i) = \beta(j)\}|,$$
$$m(\alpha, \beta) = |\{i \in \{1, 2, 3\} : \alpha(i) = \beta(i)\}|,$$

where $|X|$ denotes the number of elements of a finite set X .

1.1. Lemma. *For any $\alpha, \beta \in C$, we have:*

$$1 + n(\alpha, \beta) \geq m(\alpha, \beta).$$

Proof. The inequality is trivial if $m(\alpha, \beta)$ is either 0 or 1. If $m(\alpha, \beta) = 2$, a simple inspection of possibilities shows that $n(\alpha, \beta) \geq m(\alpha, \beta)$. Finally, if $m(\alpha, \beta) = 3$ then $\alpha = \beta$ and $n(\alpha, \beta) = n(\alpha, \alpha) \geq 2$. \square

By a *probability distribution* on C we mean a map $p : C \rightarrow [0, 1]$ such that:

$$\sum_{\alpha \in C} p(\alpha) = 1.$$

We denote by Δ the set of all probability distributions on C . Given $p, q \in \Delta$, $i, j \in \{1, 2, 3\}$, we set:

$$(1.1) \quad P_{ij}(p, q) = \sum_{\substack{\alpha, \beta \in C \\ \alpha(i) = \beta(j)}} p(\alpha)q(\beta).$$

1.2. Lemma. *Given $p, q \in \Delta$ then:*

$$(1.2) \quad 1 + \sum_{\substack{i, j=1 \\ i \neq j}}^3 P_{ij}(p, q) \geq \sum_{i=1}^3 P_{ii}(p, q).$$

Proof. We have:

$$\sum_{\substack{i, j=1 \\ i \neq j}}^3 P_{ij}(p, q) = \sum_{\substack{i, j=1 \\ i \neq j}}^3 \sum_{\substack{\alpha, \beta \in C \\ \alpha(i) = \beta(j)}} p(\alpha)q(\beta) = \sum_{\substack{i, j=1 \\ i \neq j}}^3 \sum_{\alpha, \beta \in C} \chi(i, j, \alpha, \beta) p(\alpha)q(\beta),$$

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where $\chi(i, j, \alpha, \beta) = 1$ for $\alpha(i) = \beta(j)$ and $\chi(i, j, \alpha, \beta) = 0$ for $\alpha(i) \neq \beta(j)$.
But:

$$\begin{aligned} \sum_{\substack{i,j=1 \\ i \neq j}}^3 \sum_{\alpha, \beta \in C} \chi(i, j, \alpha, \beta) p(\alpha) q(\beta) &= \sum_{\alpha, \beta \in C} \sum_{\substack{i,j=1 \\ i \neq j}}^3 \chi(i, j, \alpha, \beta) p(\alpha) q(\beta) \\ &= \sum_{\alpha, \beta \in C} n(\alpha, \beta) p(\alpha) q(\beta), \end{aligned}$$

and therefore:

$$(1.3) \quad \sum_{\substack{i,j=1 \\ i \neq j}}^3 P_{ij}(p, q) = \sum_{\alpha, \beta \in C} n(\alpha, \beta) p(\alpha) q(\beta).$$

A similar computation gives:

$$(1.4) \quad \sum_{i=1}^3 P_{ii}(p, q) = \sum_{i=1}^3 \sum_{\substack{\alpha, \beta \in C \\ \alpha(i) = \beta(i)}} p(\alpha) q(\beta) = \sum_{\alpha, \beta \in C} m(\alpha, \beta) p(\alpha) q(\beta).$$

Hence:

$$\begin{aligned} 1 + \sum_{\substack{i,j=1 \\ i \neq j}}^3 P_{ij}(p, q) &\stackrel{(1.3)}{=} 1 + \sum_{\alpha, \beta \in C} n(\alpha, \beta) p(\alpha) q(\beta) \\ &= \sum_{\alpha, \beta \in C} p(\alpha) q(\beta) + \sum_{\alpha, \beta \in C} n(\alpha, \beta) p(\alpha) q(\beta) \\ &= \sum_{\alpha, \beta \in C} (1 + n(\alpha, \beta)) p(\alpha) q(\beta) \geq \sum_{\alpha, \beta \in C} m(\alpha, \beta) p(\alpha) q(\beta) \stackrel{(1.4)}{=} \sum_{i=1}^3 P_{ii}(p, q), \end{aligned}$$

where the last inequality follows from Lemma 1.1. \square

Now let μ be an arbitrary probability measure¹ on the cartesian product $\Delta^2 = \Delta \times \Delta$. We set:

$$(1.5) \quad \mathcal{P}_{ij}(\mu) = \int_{\Delta^2} P_{ij}(p, q) d\mu(p, q),$$

for all $i, j \in \{1, 2, 3\}$.

¹We assume that the domain of μ is a σ -algebra of subsets of $\Delta \times \Delta$ that is large enough to make the map $(p, q) \mapsto p(\alpha)q(\beta)$ measurable, for all $\alpha, \beta \in C$. Notice that Δ can be identified with a subset of \mathbb{R}^3 (an affine simplex) and that such condition holds if we endow $\Delta \times \Delta$ with the Borel σ -algebra.

1.3. Proposition. *Given a probability measure μ on Δ^2 , the following inequality holds:*

$$1 + \sum_{\substack{i,j=1 \\ i \neq j}}^3 \mathcal{P}_{ij}(\mu) \geq \sum_{i=1}^3 \mathcal{P}_{ii}(\mu).$$

Proof. Follows immediately by taking the integral with respect to μ on both sides of inequality (1.2). \square

2. AN EXPERIMENT

Assume that we have a pair of boxes, call them A and B . Each box has three buttons, numbered 1, 2, 3 and two small lamps; one lamp is colored green and the other is colored red. If we press one of the buttons in a box then one (and only one) of the colored lamps lights up. The two boxes are manufactured by the same source (let's call it the *factory*); then, box A is sent to a far away location and box B is sent to another far away location. Each box will be received by an experimenter that will randomly choose one of the three buttons and press it. The experiment is repeated a large number of times, each time with a new pair of boxes. Each time the pair of boxes is prepared, the factory chooses a pair $p, q \in \Delta$ of probability distributions on the set C of all possible triples of colors. The probability distribution p is programmed into the box A as follows: when button number $i \in \{1, 2, 3\}$ is pressed, the box randomly chooses a triple of colors $\alpha \in C$ according to the probability distribution p and then it lights up the lamp² colored $\alpha(i)$. Similarly, the probability distribution q is programmed into the box B . The factory chooses the pairs (p, q) of probability distributions randomly according to some probability measure μ on $\Delta^2 = \Delta \times \Delta$. It seems that this procedure is general enough so that essentially any manufacturing strategy used by the factory is covered, *as long as the boxes A and B are not allowed to interact after a button is pressed by one of the experimenters*. This “impossibility of interaction” hypothesis — let's call it *locality* — is encoded in our formalism in the hypothesis that the lottery p used by box A to associate a button to a lamp is independent of the lottery q used by box B to associate a button to a lamp. Notice that, since p and q are chosen by the same source, we allow the choice of q to be dependent on the choice of p ; this is encoded in our formalism in the fact that the probability measure μ is *arbitrary*, i.e., it is not necessarily a product $\mu_1 \times \mu_2$ of two probability measures μ_1, μ_2 on Δ .

Now, let $p, q \in \Delta$ be given and assume that the boxes A and B are programmed respectively with the probability distributions p and q . Given, $i, j \in \{1, 2, 3\}$, if button number i is pressed in box A and button number j

²The factory could choose a triple $p_1, p_2, p_3 \in [0, 1]$ and programme box A to behave as follows: if button number $i \in \{1, 2, 3\}$ is pressed then the green lamp lights up with probability p_i (and the red lamp lights up with probability $1 - p_i$). This would amount to defining $p \in \Delta$ by the formula: $p(\alpha) = \prod_{i=1}^3 (p_i + \alpha(i)(1 - 2p_i))$, for all $\alpha \in C$.

is pressed in box B , the probability that both boxes will light up the same color is $P_{ij}(p, q)$ (recall (1.1)). Now recall that the experiment is repeated a large number N of times. Given $i, j \in \{1, 2, 3\}$, let N_{ij} denote the number of times that button number i was pressed in box A and button number j was pressed in box B . Let N_{ij}^* denote the number of times that button number i was pressed in box A , button number j was pressed in box B and the same color lit up in both boxes. We now assume that the experimenters choose the buttons they press using a lottery that is independent³ of the lottery μ used by the factory to choose the pair (p, q) . This hypothesis implies that the quotient of N_{ij}^* by N_{ij} is approximately equal to $\mathcal{P}_{ij}(\mu)$ (recall (1.5)); such approximation gets better as N gets larger.

Now, assume that the experiment has been performed a large number N of times and that we have observed:

$$(2.1) \quad \frac{N_{ij}^*}{N_{ij}} \approx \frac{1}{4}, \quad \text{for } i, j \in \{1, 2, 3\}, i \neq j,$$

$$(2.2) \quad \frac{N_{ii}^*}{N_{ii}} \approx 1, \quad \text{for } i \in \{1, 2, 3\},$$

where \approx denotes⁴ “approximately equal”. Is this possible? We have:

$$\mathcal{P}_{ij}(\mu) \approx \frac{1}{4}, \quad \text{for } i \neq j, \quad \mathcal{P}_{ii}(\mu) \approx 1,$$

so that:

$$1 + \sum_{\substack{i,j=1 \\ i \neq j}}^3 \mathcal{P}_{ij}(\mu) \approx 1 + \frac{6}{4} = \frac{5}{2},$$

$$\sum_{i=1}^3 \mathcal{P}_{ii}(\mu) \approx 3.$$

But this contradicts Proposition 1.3.

Now, it is a well-known fact that such an experiment can be performed⁵ and that both (2.1) and (2.2) are observed. What is going on? Since these observations contradict Proposition 1.3, we have to drop one of our hypotheses: either the “non clairvoyant factory hypothesis” or the “locality hypothesis”. It seems that locality is gone for good!

³We call this the *non clairvoyant factory* hypothesis: the factory cannot predict how the buttons to be pressed (far away in space and in the future) will be chosen.

⁴How good the approximation must be? Good enough so that I can get to the contradiction with Proposition 1.3 explained below.

⁵See [1, 2, 3]. The experiments confirm the quantum predictions for measurements of polarization of entangled pairs of photons. The theoretical prediction gives $N_{ii}^* = N_{ii}$ and $\mathcal{P}_{ii}(\mu) = 1$, but we can never count on real life experiments to give us perfect correlations, so my argument allows $\mathcal{P}_{ii}(\mu) \approx 1$.

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