

A Proof for a Conjecture of Gorgol

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1 Introduction

The *Turán number* $ex(n, F)$ is the maximum number of edges in a graph on n vertices which does not contain H as a subgraph. For a general graph F , the value of $ex(n, F)$ is asymptotically well known and dependent of the *chromatic number* $\chi(F)$. Erdős, Stone and Simonovits' Theorem [1] state that

$$\lim_{n \rightarrow \infty} \frac{ex(n, F)}{n^2} = \frac{\chi(F) - 2}{2\chi(F) - 2}.$$

Although this theorem gives a lot of information on the asymptotic growth of $ex(n, F)$, it should be noted that it is only of interest for nonbipartite graphs. If F is bipartite, it asserts merely that $ex(n, F) = o(n^2)$. Therefore, we focus on the Turán number for bipartite graphs. In particular, we consider when $F = kP_3$, where P_r is a path on r vertices and kG consists of k disjoint copies of the graph G .

For graphs G and F , let $G + F$ and $G \vee F$ be the *disjoint union* and *join* of G and F , respectively. The join of G and F is the graph obtained from $G + F$ by adding edges from all vertices of G to all vertices of F . Let $H_{ex}(n, G)$ represent an *extremal graph* on n vertices without G as subgraph with $ex(n, G)$ edges.

Gorgol [2] gave upper and lower bounds for the Turán number forbidding kF , for a graph F on r vertices. The lower bound was obtained by noting that neither

$$\begin{aligned} G_1(n, kF) &= K_{kr-1} + H_{ex}(n - kr + 1, F) \\ G_2(n, kF) &= K_{k-1} \vee H_{ex}(n - k + 1, F) \end{aligned}$$

contain k disjoint copies of F . Indeed, $G_1(n, kF)$ contains only $k - 1$ copies of F in K_{kr-1} and any copy of F in $G_2(n, kF)$ must contain at least one vertex in K_{k-1} . Gorgol's lower bound is, therefore,

$$ex(n, kF) \geq \max\{e(G_1(n, kF)), e(G_2(n, kF))\}$$

where $e(G)$ denotes the number of edges in a graph G .

2 Disjoint copies of P_3

We consider the case when $F = P_3$. Let

$$\begin{aligned} g_1(n, k) &= e(G_1(n, kP_3)), \\ g_2(n, k) &= e(G_2(n, kP_3)) \text{ and} \\ \text{Gorgol}(n, k) &= \max\{g_1(n, k), g_2(n, k)\}. \end{aligned}$$

Note that $H_{ex}(n, P_3) = M_n$, where M_n is a *near perfect matching* on n vertices. If n is even, then M_n consists of a matching on n vertices and, if n is odd, then $M_n = M_{n-1} + K_1$. Thus, $e(M_n) = \lfloor n/2 \rfloor$ and we have

$$\begin{aligned} g_1(n, k) &= \binom{3k-1}{2} + \left\lfloor \frac{n-3k+1}{2} \right\rfloor; \\ g_2(n, k) &= \binom{k-1}{2} + (k-1)(n-k+1) + \left\lfloor \frac{n-k+1}{2} \right\rfloor; \\ \text{Gorgol}(n, k) &= \begin{cases} g_1(n, k), & \text{for } 3k \leq n \leq 5k, \\ g_2(n, k), & \text{for } n \geq 5k. \end{cases} \end{aligned}$$

Gorgol's Conjecture states that, when $n \geq 3k$,

$$ex(n, kP_3) = \text{Gorgol}(n, K).$$

In the same paper, Gorgol proved this is true when $k \in \{2, 3\}$. Bushaw and Kettle [3] also proved the conjecture is true for $n \geq 7k$.

In this work, we give a constructive proof for Gorgol's Conjecture for all values of n and k . In particular, we give an algorithm that finds a set of disjoint copies of P_3 given a graph G as input. We prove Gorgol's Conjecture by showing that, if G has n vertices and more than $\text{Gorgol}(n, k)$ edges, then this algorithm returns at least k disjoint copies of P_3 in G .

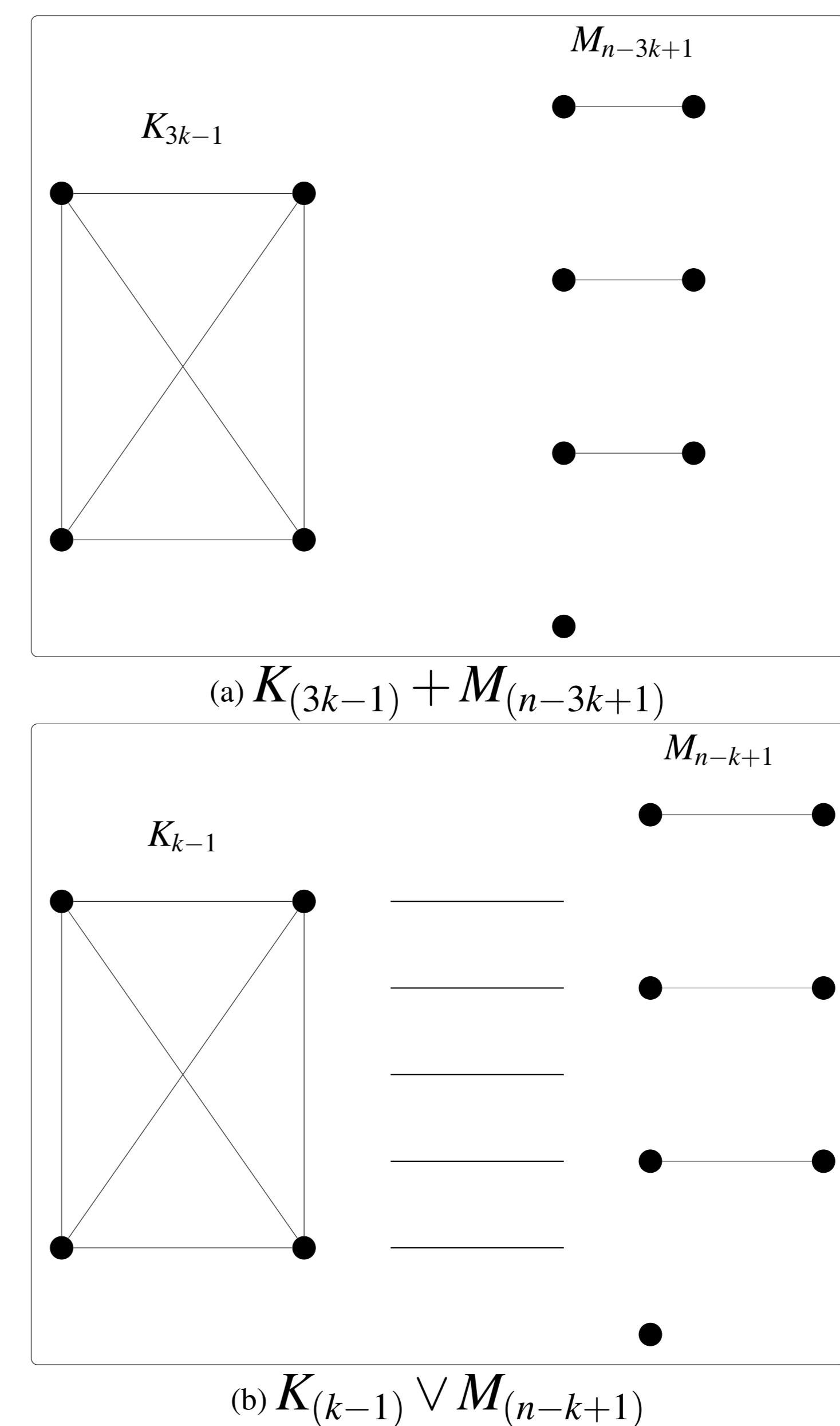
3 Algorithm

The algorithm receives as input a graph $G = (V, E)$ and an integer $k \leq n/3$.

Let $\mathcal{Q} = \{Q_1, Q_2, \dots, Q_s\}$ be a collection of subsets of $V(G)$ such that:

- $Q_i \cap Q_j = \emptyset$.
- $|Q_i| = 3$.
- $G[Q_i]$ contains a copy of P_3 as subgraph.

We start with an empty collection \mathcal{Q} and iteratively look for an improvement \mathcal{Q}' for \mathcal{Q} . We say that a collection \mathcal{Q}' is an improvement of \mathcal{Q} if either $|\mathcal{Q}'| > |\mathcal{Q}|$ or $|\mathcal{Q}'| = |\mathcal{Q}|$ and \mathcal{Q}' has more triangles than \mathcal{Q} .



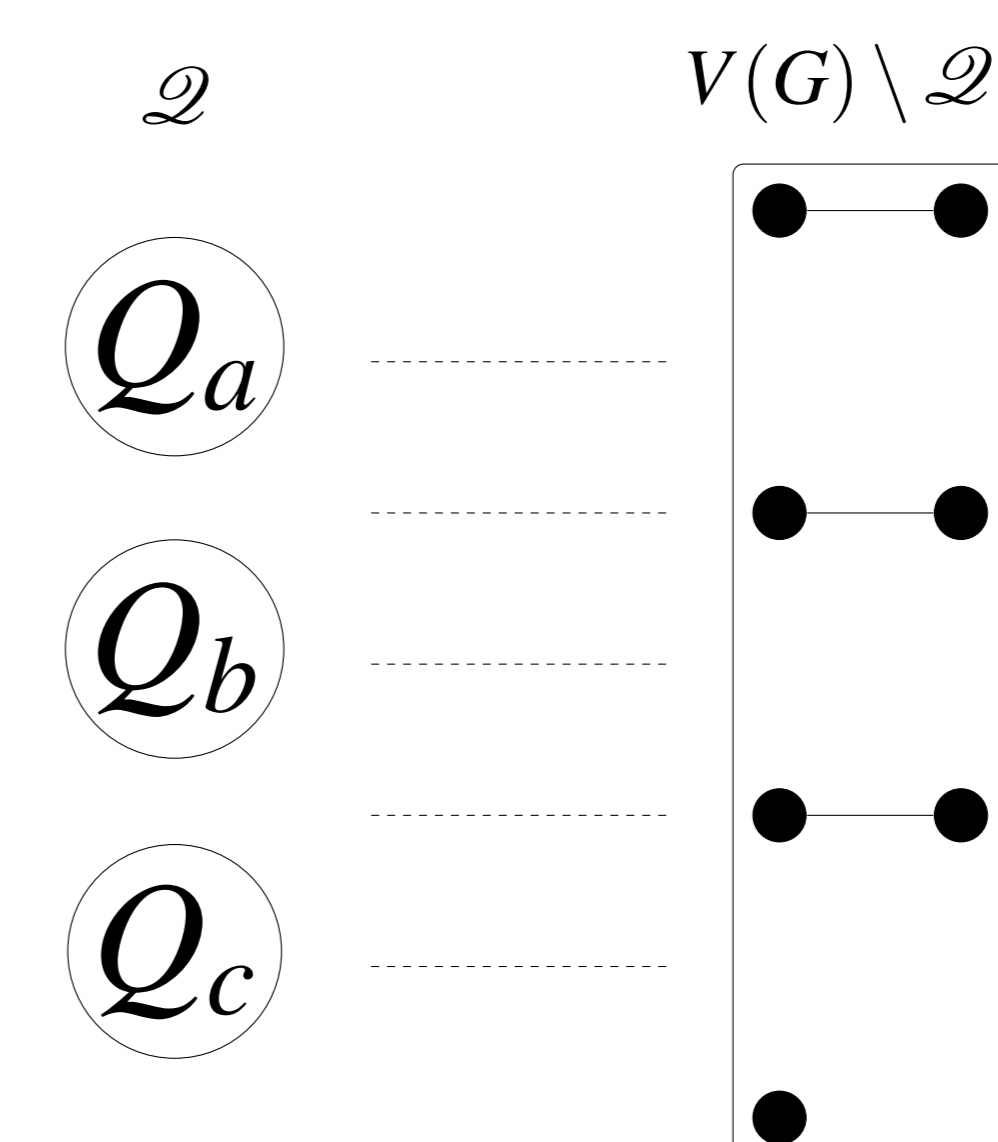
The given algorithm is then as follows:

Algorithm 1 Disjoint copies of P_3 .

Input: (G, k) .

- 1: $\mathcal{Q} \leftarrow \emptyset$
- 2: **while** \exists improvement \mathcal{Q}' of \mathcal{Q} **do**
- 3: $\mathcal{Q} \leftarrow \mathcal{Q}'$
- 4: **end while**
- 5: **return** \mathcal{Q}

To find improvements for our current collection \mathcal{Q} , we need to search only in a small subgraph of G .



With local improvements only we can find the desired k disjoint copies of P_3 , proving the conjecture. The naive complexity of the algorithm is $O(n^{12}k^5)$. With a subset of local improvements and good use of data structures, we find a collection of k disjoint P_3 in $O(k|E|)$ time. Our main result, which proves Gorgol's Conjecture, can then be stated as

Theorem. *If a graph G has n vertices and $e(G) \geq \text{Gorgol}(n, k)$ edges, then the provided algorithm will find a collection \mathcal{Q} of size at least k .*

References

- [1] Bondy, J. A. and Murty, U.S.R. *Graph Theory* Graduate Texts in Mathematics 244, 2008, pp.317-320
- [2] Gorgol, I. *Turán Numbers fo disjoint copies of graphs*. Graphs and Combinatorics (4 January 2011), pp. 1-7.
- [3] Bushaw, N. and Kettle, N. *Turán Numbers of multiple paths and equibipartite forests*. Combinatorics, Probability and Computing, 2011, Vol.20(6), pp.837-853