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Simpler, Faster, Better: Concurrency Utilities in JDK™ Software Version 5.0

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Goal

Learn how to use the new concurrency utilities (the `java.util.concurrent` package) to replace error-prone or inefficient code and to better structure applications

Agenda

Overview of `java.util.concurrent`

Concurrent Collections

Threads Pools and Task Scheduling

Locks, Conditions and Synchronizers

Atomic Variables

Rationale for `java.util.concurrent`

Developing Concurrent Classes Was Just Too Hard

- The built-in concurrency primitives—`wait()`, `notify()`, and `synchronized`—
Are, well, primitive
 - Hard to use correctly
 - Easy to use incorrectly
 - Specified at too low a level for most applications
 - Can lead to poor performance if used incorrectly
- Too much wheel-reinventing!

Goals for `java.util.concurrent`

Simplify Development of Concurrent Applications

- Provide a set of basic concurrency building blocks
- Something for everyone
 - Make some problems trivial to solve by everyone
 - Develop thread-safe classes, such as servlets, built on concurrent building blocks like `ConcurrentHashMap`
 - Make some problems easier to solve by concurrent programmers
 - Develop concurrent applications using thread pools, barriers, latches, and blocking queues
 - Make some problems possible to solve by concurrency experts
 - Develop custom locking classes, lock-free algorithms

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Concurrent Collections

Concurrent vs. Synchronized

- Pre Java™ 5 platform: **thread-safe but not concurrent classes**
- Thread-safe synchronized collections
 - `Hashtable`, `Vector`, `Collections.synchronizedMap`
 - Monitor is source of contention under concurrent access
 - Often require locking during iteration
- Concurrent collections
 - Allow multiple operations to overlap each other
 - Big performance advantage
 - At the cost of some slight differences in semantics
 - Might not support atomic operations

Concurrent Collections

- **ConcurrentHashMap**
 - Concurrent (scalable) replacement for **Hashtable** or **Collections.synchronizedMap**
 - Allows reads to overlap each other
 - Allows reads to overlap writes
 - Allows up to 16 writes to overlap
 - Iterators don't throw **ConcurrentModificationException**
- **CopyOnWriteArrayList**
 - Optimized for case where iteration is much more frequent than insertion or removal
 - Ideal for event listeners

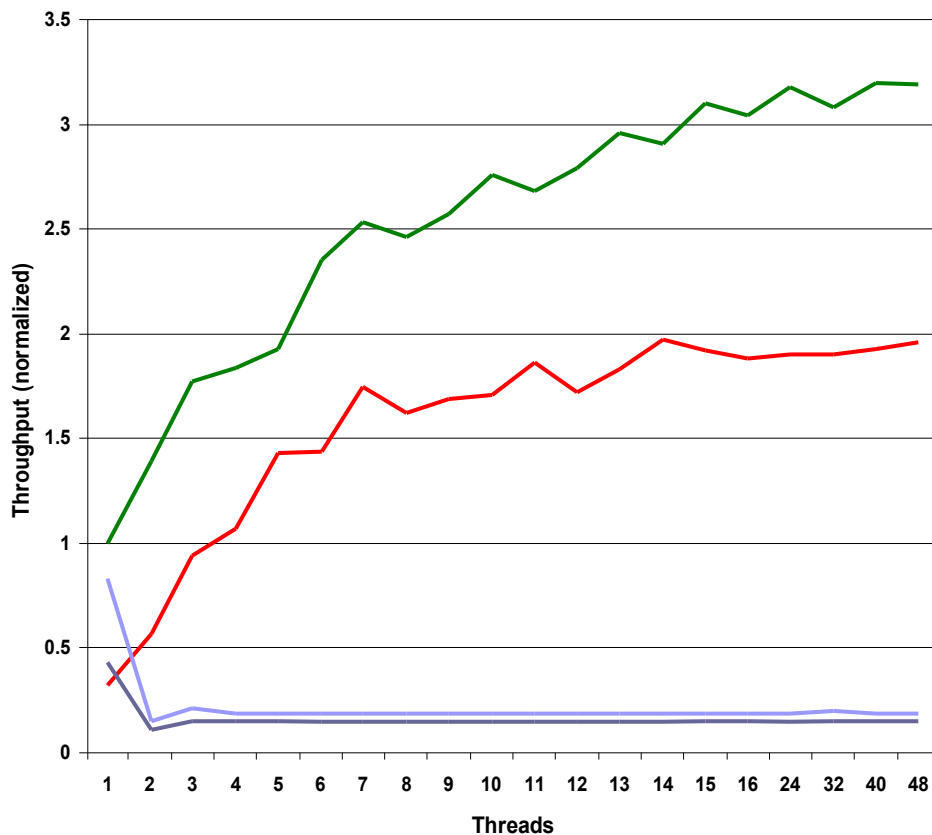
Concurrent Collections

Iteration Semantics

- Synchronized collection iteration broken by concurrent changes in another thread
 - Throws **ConcurrentModificationException**
 - Locking a collection during iteration hurts scalability
- Concurrent collections can be modified concurrently during iteration
 - Without locking the whole collection
 - Without **ConcurrentModificationException**
 - But changes may not be seen

Concurrent Collection Performance

Throughput in Thread-safe Maps



- ConcurrentHashMap
- ConcurrentSkipListMap
- SynchronizedHashMap
- SynchronizedTreeMap

Java 6 B77
8-way system
40% read only
60% insert
2% removals

Should **concurrent...** be courier new?

Queues

New Interface Added to `java.util`

```
interface Queue<E> extends Collection<E> {
    boolean offer(E x);
    E poll();
    E remove() throws NoSuchElementException;
    E peek();
    E element() throws NoSuchElementException;
}
```

- Retrofit (non-thread-safe)–implemented by **LinkedList**
- Add (non-thread-safe) **PriorityQueue**
- Fast thread-safe non-blocking **ConcurrentLinkedQueue**
- Better performance than **LinkedList** is possible as random-access requirement has been removed

Blocking Queues

BlockingQueueInterface

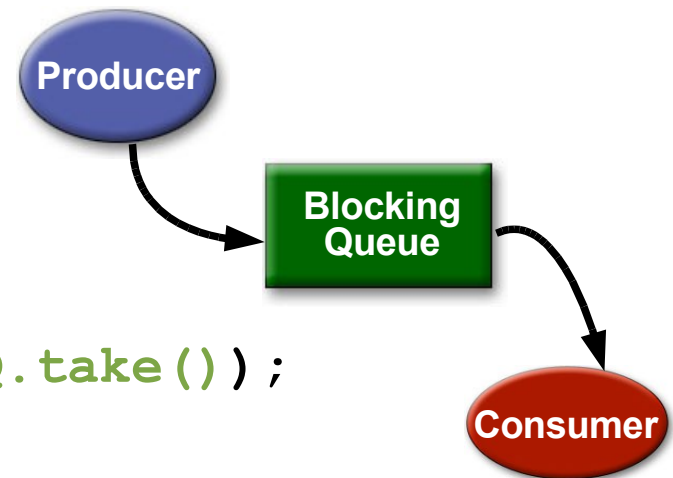
- Extends **Queue** to provides blocking operations
 - Retrieval: **take**—Wait for queue to become nonempty
 - Insertion: **put**—Wait for capacity to become available
- Several implementations:
 - **LinkedBlockingQueue**
 - Ordered FIFO, may be bounded, two-lock algorithm
 - **PriorityBlockingQueue**
 - Unordered but retrieves least element, unbounded, lock-based
 - **ArrayBlockingQueue**
 - Ordered FIFO, bounded, lock-based
 - **SynchronousQueue**
 - Rendezvous channel, lock-based in Java 5 platform, lock-free in Java 6 platform

BlockingQueue Example

```
class LogWriter {
    final BlockingQueue msgQ =
        new LinkedBlockingQueue();

    public void writeMessage(String msg) throws IE {
        msgQ.put(msg);
    }

    // run in background thread
    public void logServer() {
        try {
            while (true) {
                System.out.println(msgQ.take());
            }
        }
        catch (InterruptedException ie) { ... }
    }
}
```



Producer-Consumer Pattern

- **LogWriter** example illustrates the *producer-consumer* pattern
 - Ubiquitous concurrency pattern, nearly always relies on some form of blocking queue
 - Decouples identification of work from doing the work
 - Simpler and more flexible
- **LogWriter** had many producers, one consumer
 - Thread pool has many producers, many consumers
- **LogWriter** moves IO from caller to log thread
 - Shorter code paths, fewer context switches, no contention for IO locks → more efficient

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Executors

Framework for Asynchronous Execution

- Standardize asynchronous invocation
 - Framework to execute **Runnable** and **Callable** tasks
- Separate submission from execution policy
 - Use **anExecutor.execute(aRunnable)**
 - Not **new Thread(aRunnable).start()**
- Cancellation and shutdown support
- Usually created via **Executors** factory class
 - Configures flexible **ThreadPoolExecutor**
 - Customize shutdown methods, before/after hooks, saturation policies, queuing

Executors

Decouple Submission From Execution Policy

```
public interface Executor {  
    void execute(Runnable command);  
}
```

- Code which submits a task doesn't have to know in what thread the task will run
 - Could run in the calling thread, in a thread pool, in a single background thread (or even in another JVM™ software!)
 - Executor implementation determines **execution policy**
 - Execution policy controls resource utilization, saturation policy, thread usage, logging, security, etc
 - Calling code need not know the execution policy

Executor and ExecutorService

ExecutorService Adds Lifecycle Management

- **ExecutorService** supports both graceful and immediate shutdown

```
public interface ExecutorService extends Executor {  
    void shutdown();  
    List<Runnable> shutdownNow();  
    boolean isShutdown();  
    boolean isTerminated();  
    boolean awaitTermination(long time, TimeUnit unit)  
        throws InterruptedException  
  
    // other convenience methods for submitting tasks  
}
```

- Many useful utility methods too

Creating Executors

Factory Methods in the Executors Class

```
public class Executors {
    static ExecutorService
        newSingleThreadedExecutor ();

    static ExecutorService
        newFixedThreadPool (int poolSize);

    static ExecutorService
        newCachedThreadPool ();

    static ScheduledExecutorService
        newScheduledThreadPool (int corePoolSize);

    // additional versions specifying ThreadFactory
    // additional utility methods
}
```

Executors Example

Web Server—Poor Resource Management

```
class UnstableWebServer {  
  
    public static void main(String[] args) {  
        ServerSocket socket = new ServerSocket(80);  
        while (true) {  
            final Socket connection = socket.accept();  
            Runnable r = new Runnable() {  
                public void run() {  
                    handleRequest(connection);  
                }  
            };  
            // Don't do this!  
            new Thread(r).start();  
        }  
    }  
}
```

Executors Example

Web Server—Better Resource Management

```
class BetterWebServer {
    Executor pool = Executors.newFixedThreadPool(7);

    public static void main(String[] args) {
        ServerSocket socket = new ServerSocket(80);
        while (true) {
            final Socket connection = socket.accept();
            Runnable r = new Runnable() {
                public void run() {
                    handleRequest(connection);
                }
            };
            pool.execute(r);
        }
    }
}
```

Saturation Policies

- An Executor which execute tasks in a thread pool
 - Can guarantee you will not run out of threads
 - Can manage thread competition for CPU resources
- There is still a risk of running out of memory
 - Tasks could queue up without bound
- Solution: Use a *bounded task queue*
 - Just so happens that JUC provides several of these...
- If queue fills up, the *saturation policy* is applied
 - Policies available: Throw, discard oldest, discard newest, or run-in-calling-thread
 - The last has the benefit of throttling the load

Futures and Callables

Representing Asynchronous Tasks

- **Callable** is functional analog of **Runnable**

```
interface Callable<V> {  
    V call() throws Exception;  
}
```

- **Future** holds result of asynchronous call, normally a **Callable**

```
interface Future<V> {  
    V get() throws InterruptedException,  
        ExecutionException;  
    V get(long timeout, TimeUnit unit) throws ...;  
    boolean cancel(boolean mayInterrupt);  
    boolean isCancelled();  
    boolean isDone();  
}
```

Futures Example

Implementing a Concurrent Cache

```
public class Cache<K, V> {
    final ConcurrentMap<K, FutureTask<V>> map =
        new ConcurrentHashMap<K, FutureTask<V>>();

    public V get(final K key) throws InterruptedException {
        FutureTask<V> f = map.get(key);
        if (f == null) {
            Callable<V> c = new Callable<V>() {
                public V call() {
                    // return value associated with key
                }
            };
            f = new FutureTask<V>(c);
            FutureTask<V> old = map.putIfAbsent(key, f);
            if (old == null)
                f.run();
            else
                f = old;
        }
        try { return f.get(); }
        catch (ExecutionException ex) { // rethrow ex.getCause() }
    }
}
```


ScheduledExecutorService

Deferred and Recurring Tasks

- **ScheduledExecutorService** can be used to:
 - Schedule a **Callable** or **Runnable** to run once with a fixed delay after submission
 - Schedule a **Runnable** to run periodically at a fixed rate
 - Schedule a **Runnable** to run periodically with a fixed delay between executions
- Submission returns a **ScheduledFutureTask** handle which can be used to cancel the task
- Like **java.util.Timer**, but supports pooling

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Locks

- Use of monitor synchronization is just fine for most applications, but it has some shortcomings
 - Single wait-set per lock
 - No way to interrupt or time-out when waiting for a lock
 - Locking must be block-structured
 - Inconvenient to acquire a variable number of locks at once
 - Advanced techniques, such as hand-over-hand locking, are not possible
- Lock objects address these limitations
 - But harder to use: Need `finally` block to ensure release
 - **So if you don't need them, stick with `synchronized`**

Framework for Flexible Locking

```

interface Lock {
    void      lock();
    void      lockInterruptibly() throws
                InterruptedException;

    boolean   tryLock();
    boolean   tryLock(long time, TimeUnit unit) throws
                InterruptedException;

    void      unlock();
    Condition newCondition() throws
                UnsupportedOperationException;
}
    
```

- High-performance implementation: **ReentrantLock**
 - Basic semantics same as use of **synchronized**
 - Condition object semantics like **wait/notify**

Simple Lock Example

- Used extensively within `java.util.concurrent`

```
final Lock lock = new ReentrantLock();  
...  
lock.lock();  
try {  
    // perform operations protected by lock  
}  
catch(Exception ex) {  
    // restore invariants & rethrow  
}  
finally {  
    lock.unlock();  
}
```

- Must manually ensure lock is released

Conditions

Monitor-like Operations for Working With Locks

- **Condition** is an abstraction of **wait/notify**

```
interface Condition {
    void    await() throws InterruptedException;
    boolean await(long time, TimeUnit unit)
                throws InterruptedException;
    long    awaitNanos(long nanosTimeout)
                throws InterruptedException;
    boolean awaitUntil(Date deadline)
                throws InterruptedException;
    void    awaitUninterruptibly();

    void    signal();
    void    signalAll();
}
```

- Timed **await** versions report reason for return

Condition Example

```
class BoundedBuffer {
    final Lock lock          = new ReentrantLock();
    final Condition notFull  = lock.newCondition();
    final Condition notEmpty = lock.newCondition();
    ...
    void put(Object x) throws InterruptedException {
        lock.lock(); try {
            while (isFull()) notFull.await();
            doPut(x);
            notEmpty.signal();
        } finally { lock.unlock(); }
    }
    Object take() throws InterruptedException {
        lock.lock(); try {
            while (isEmpty()) notEmpty.await();
            notFull.signal();
            return doTake();
        } finally { lock.unlock(); }
    }
}
```

Synchronizers

Utility Classes for Coordinating Access and Control

- **Semaphore**—Dijkstra counting semaphore, managing a specified number of permits
- **CountDownLatch**—allows one or more threads to wait for a set of threads to complete an action
- **CyclicBarrier**—allows a set of threads to wait until they all reach a specified barrier point
- **Exchanger**—allows two threads to rendezvous and exchange data
 - Such as exchanging an empty buffer for a full one

Semaphore Example

Bound the Submission of Tasks to an Executor

```
public class ExecutorProxy implements Executor {
    private final Semaphore tasks;
    private final Executor master;

    ExecutorProxy(Executor master, int limit) {
        this.master = master;
        tasks = new Semaphore(limit);
    }

    public void execute(Runnable r) {
        tasks.acquireUninterruptibly(); // for simplicity
        try {
            master.execute(r);
        }
        finally {
            tasks.release();
        }
    }
}
```

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Atomic Variables

Holder Classes for Scalars, References and Fields

- Support atomic operations
 - Compare-and-set (CAS)
 - Get, set and arithmetic operations (where applicable)
 - Increment, decrement operations
- Abstraction of **volatile** variables
- Nine main classes:
 - { **int**, **long**, **reference** } X { **value**, **field**, **array** }
- e.g. **AtomicInteger** useful for counters, sequence numbers, statistics gathering

AtomicInteger Example

Construction Counter for Monitoring/Management

- Replace this:


```
class Service {
    static int services;
    public Service() {
        synchronized (Service.class) {
            services++;
        }
    } // ...
}
```

- With this:


```
class Service {
    static AtomicInteger services =
        new AtomicInteger();
    public Service() {
        services.getAndIncrement();
    }
    // ...
}
```

Atomic Compare-and-Set (CAS)

- `boolean compareAndSet(int expected, int update)`
 - Atomically sets value to `update` if currently `expected`
 - Returns true on successful update
- Direct hardware support in all modern processors
 - CAS, cmpxchg, ll/sc
- High-performance on multi-processors
 - No locks, so no lock contention and no blocking
 - But can fail
 - So algorithms must implement retry loop
- Foundation of many concurrent algorithms

Sneak Preview of Java 6 Platform (Code-Named Mustang)

- Double-ended queues: `Deque`, `BlockingDeque`
 - Implementations: `ArrayDeque`, `LinkedBlockingDeque`, `ConcurrentLinkedDeque`
- Concurrent skiplists: `ConcurrentSkipList{Map|Set}`
- Enhanced navigation of sorted maps/sets
 - `Navigable{Map|Set}`
- Miscellaneous algorithmic enhancements
 - More use of lock-free algorithms in utilities
 - VM performance improvements for intrinsic locking
- M&M support for locks and conditions

java.util.concurrent

- Executors
 - Executor
 - ExecutorService
 - ScheduledExecutorService
 - Callable
 - Future
 - ScheduledFuture
 - Delayed
 - CompletionService
 - ThreadPoolExecutor
 - ScheduledThreadPoolExecutor
 - AbstractExecutorService
 - Executors
 - FutureTask
 - ExecutorCompletionService
- Queues
 - BlockingQueue
 - ConcurrentLinkedQueue
 - LinkedBlockingQueue
 - ArrayBlockingQueue
 - SynchronousQueue
 - PriorityBlockingQueue
 - DelayQueue
- Concurrent collections
 - ConcurrentHashMap
 - ConcurrentHashMap
 - CopyOnWriteArray{List, Set}
- Synchronizers
 - CountdownLatch
 - Semaphore
 - Exchanger
 - CyclicBarrier
- Locks: java.util.concurrent.locks
 - Lock
 - Condition
 - ReadWriteLock
 - AbstractQueuedSynchronizer
 - LockSupport
 - ReentrantLock
 - ReentrantReadWriteLock
- Atomics: java.util.concurrent.atomic
 - Atomic[Type]
 - Atomic[Type]Array
 - Atomic[Type]FieldUpdater
 - Atomic{Markable, Stampable}Reference

Summary

- Whenever you are about to use
 - `Object.wait`, `notify`, `notifyAll`
 - `new Thread(aRunnable).start()` ;
 - `synchronized`
- Check first in `java.util.concurrent` if there is a class that ...
 - Does it already, or
 - Let's you do it a simpler, or better way, or
 - Provides a better starting point for your own solution
- Don't reinvent the wheel!

For More Information

- Javadoc™ tool for `java.util.concurrent`—
In JDK™ 5.0 software download or on Sun website
- Doug Lea’s concurrency-interest mailing list
 - <http://gee.cs.oswego.edu/dl/concurrency-interest/index.html>
- **Concurrent Programming in Java** (Lea)
 - Addison-Wesley, 1999 ISBN 0-201-31009-0
- **Java Concurrency in Practice** (Goetz, et al)
 - Addison-Wesley, 2006, ISBN 0-321-34960-1
- JUC Backport to JDK 1.4 software
 - <http://www.mathcs.emory.edu/dcl/util/backport-util-concurrent/>

Q&A

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